

AN INTEGRATED FIRST-ORDER THEORY OF POINTS AND INTERVALS OVER LINEAR ORDERS (PART II)

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ABSTRACT. There are two natural and well-studied approaches to temporal ontology and reasoning: point-based and interval-based. Usually, interval-based temporal reasoning deals with points as a particular case of duration-less intervals. A recent result by Balbiani, Goranko, and Sciavicco presented an explicit two-sorted point-interval temporal framework in which time instants (points) and time periods (intervals) are considered on a par, allowing the perspective to shift between these within the formal discourse. We consider here two-sorted first-order languages based on the same principle, and therefore including relations, as first studied by Reich, among others, between points, between intervals, and inter-sort. We give complete classifications of its sub-languages in terms of relative expressive power, thus determining how many, and which, are the intrinsically different extensions of two-sorted first-order logic with one or more such relations. This approach roots out the classical problem of whether or not points should be included in an interval-based semantics. In this Part II, we deal with the cases of all dense and the case of all unbounded linearly ordered sets.

1. INTRODUCTION

The relevance of temporal logics in many theoretical and applied areas of computer science and AI, such as theories of action and change, natural language analysis and processing, and constraint satisfaction problems, is widely recognized. While the predominant approach in the study of temporal reasoning and logics has been based on the assumption that time

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points (instants) are the primary temporal ontological entities, there has also been significant activity in the study of interval-based temporal reasoning and logics over the past two decades. The variety of binary relations between intervals in linear orders was first studied systematically by Allen [AH87, All83, AF94], who explored their use in systems for time management and planning. Allen’s work and much that follows from it is based on the assumption that time can be represented as a dense line, and that points are excluded from the semantics. At the modal level, Halpern and Shoham [HS91] introduced the multi-modal logic HS that comprises modal operators for all possible relations (known as Allen’s relations [All83]) between two intervals in a linear order, and it has been followed by a series of publications studying the expressiveness and decidability/undecidability and complexity of the fragments of HS, e.g., [BMM⁺14, BMG⁺14]. Many studies on interval logics have considered the so-called ‘non-strict’ interval semantics, allowing point-intervals (with coinciding endpoints) along with proper ones, and thus encompassing the instant-based approach, too; more recent ones, instead, started to treat pure intervals only. Yet, little has been done so far on the formal treatment of both temporal primitives, points and intervals, in a unified two-sorted framework. A detailed philosophical study of both approaches, point-based and interval-based, can be found in [vB91] (see also [CM00]). A similar mixed approach has been studied in [AH89]. [MH06] contains a study of the two sorts and the relations between them in dense linear orders. More recently, a modal logic that includes different operators for points and interval has been presented in [BGS11].

The present paper provides a systematic treatment of point and interval relations (including equality between points and between intervals treated on the same footing as the other relations) at the first-order level. Our work is motivated, among other observations, by the fact that natural languages incorporate both ontologies on a par, without assuming the primacy of one over the other, and have the capacity to shift the perspective smoothly from instants to intervals and vice versa within the same discourse, e.g.: *when the alarm goes on, it stays on until the code is entered*, which contains two instantaneous events and a non-instantaneous one. Moreover, there are various temporal scenarios which neither of the two ontologies alone can grasp properly since neither the treatment of intervals as the sets of their internal points, nor the treatment of points as ‘instantaneous’ intervals, is really adequate. The technical identification of intervals with sets of their internal points, or of points as instantaneous intervals leads also to conceptual problems like the confusion of events and fluents. Instantaneous events are represented by time intervals and should be distinguished from instantaneous holding of fluents, which are evaluated at time points: therefore, the point a should be distinguished from the interval $[a, a]$, and the truths in these should not necessarily imply each other. Finally, we note that, while differences in expressiveness have been found between the strict and non-strict semantics for some interval logics (see [MGMS11], for example), so far, no distinction in the decidability of the satisfiability has been found. Therefore, we believe that an attempt to systemize the role of points, intervals, and their interaction, would make good sense not only from a purely ontological point of view, but also from algorithmic and computational perspectives.

Previous Work and Motivations. As presented in the early work of van Benthem [vB91] and Allen and Hayes [AH85], interval temporal reasoning can be formalized as an extension of first-order logic with equality with one or more relations, and the properties of the resulting language can be studied; obviously, the same applies when relations between points are considered too. In this paper we ask the question: interpreted over linear orders, how

many and which expressively different languages can be obtained by enriching first-order logic with relations between intervals, between points, and between intervals and points? Since, as we shall see, there are 26 different relations (including equality of both sorts) between points, intervals, and points and intervals, 2^{26} is an upper bound on this number. (It is worth noticing that in [MH06] the authors distinguish 30 relations, instead of 26; this is due to the fact that the concepts of the point *a starting* the interval $[a, b]$ and *meeting* it are considered to be different.) However, since certain relations are definable in terms of other ones, the actual number is less and in fact, as we shall show, much less. The answer also depends on our choices of certain semantic parameters, specifically, the class of linear orders over which we construct our interval structures. In this paper, in Part I [CDS18], we consider the classification problem relative to:

- (i) the class of all linear orders;
- (ii) the class of all weakly discrete linear orders (i.e., orders in which every point with a successor/predecessor has an immediate one).

In Part II of this paper we consider:

- (iii) the class of all dense linear orders;
- (iv) the class of all unbounded linear orders;

Apart from the intrinsic interest and naturalness of this classification problem, its outcome has some important repercussions, principally in the reduction of the number of cases that need to be considered in other problems relating to these languages. For example, it reduces the number of representation theorems that are needed: given the *dual* nature of time intervals (i.e., they can be abstract first-order individuals with specific characteristics, or they can be defined as ordered pairs over a linear order), one of the most important problems that arises is the existence or not of a *representation theorem*. Consider any class of linear orders: given a specific extension of first-order logic with a set of interval relations (such as, for example, *meets* and *during*), does there exist a set of axioms in this language which would constrain (abstract) models of this signature to be isomorphic to concrete ones? Various representation theorems exist in the literature for languages that include interval relations only: van Benthem [vB91], over rationals and with the interval relations *during* and *before*, Allen and Hayes [AH85], for the dense unbounded case without point intervals and for the relation *meets*, Ladkin [Lad78], for point-based structures with a quaternary relation that encodes meeting of two intervals, Venema [Ven91], for structures with the relations *starts* and *finishes*, Goranko, Montanari, and Sciavicco [GMS03], for linear structures with *meets* and *met-by*, Bochman [Boc90], for point-interval structures, and Coetzee [Coe09] for dense structure with *overlaps* and *meets*. Clearly, if two sets of relations give rise to expressively equivalent languages, two separate representation theorems for them are not needed. In which cases are representation theorems still outstanding? Preliminary works that provide similar classifications appeared in [CS11] for first-order languages with equality and only interval-interval relations, and in [CDS12] for points and intervals (with equality between intervals treated on a par with the other relations) but only over the class of all linear orders. Finally, a complete study of first-order interval temporal logics enables a deeper understanding of their modal counterparts based on their shared relational semantics.

Structure of the paper. This paper is structured as follows. Section 2 provides the necessary preliminaries, along with an overview of the general methodology used in this paper. Part I of this paper dealt with definability and undefinability in the classes *Lin* and *Dis*, from which we start in order to tackle, in Section 3, the study the expressive power of

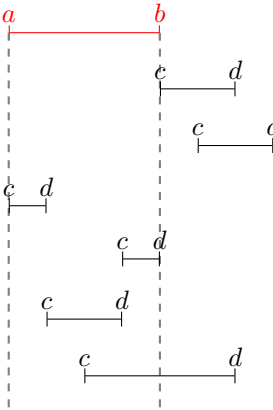
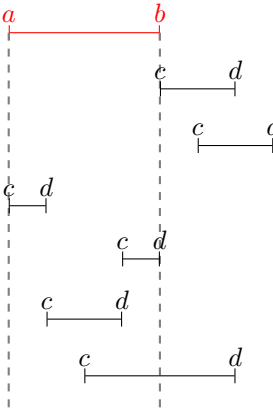
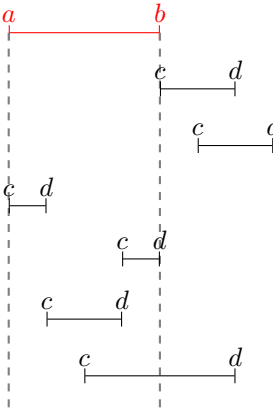
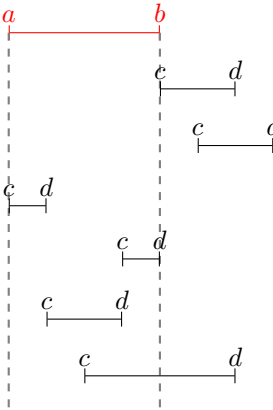
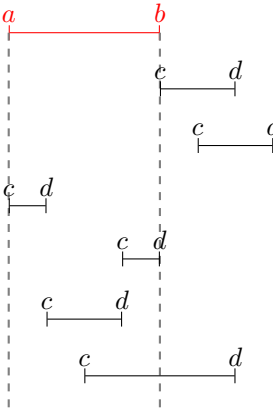
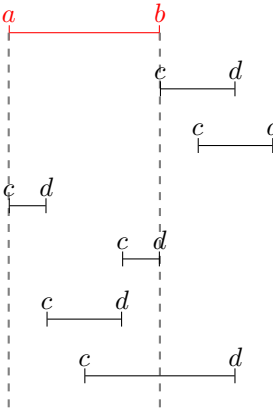
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|-----------------|----------|-----------|--|--|
| $(meets, m)$ | $[a, b]$ | 34_{ii} | $[c, d] \Leftrightarrow b = c$ |  |
| $(before, b)$ | $[a, b]$ | 44_{ii} | $[c, d] \Leftrightarrow b < c$ |  |
| $(starts, s)$ | $[c, d]$ | 14_{ii} | $[a, b] \Leftrightarrow a = c, d < b$ |  |
| $(finishes, f)$ | $[c, d]$ | 03_{ii} | $[a, b] \Leftrightarrow b = d, a < c$ |  |
| $(during, d)$ | $[c, d]$ | 04_{ii} | $[a, b] \Leftrightarrow a < c, d < b$ |  |
| $(overlaps, o)$ | $[a, b]$ | 24_{ii} | $[c, d] \Leftrightarrow a < c < b < d$ |  |

Table 1: Interval-interval relations, a.k.a. Allen’s relations. The equality relation is not depicted.

the language by analyzing the definability properties of each basic relation in the class Den , and in Section 4 the corresponding undefinability results in this case. Then, in Section 5 and Section 6, respectively, we present the same analysis in the unbounded case, before concluding. It is worth reminding that most of the results presented here are a consequence of those presented in Part I, to which we shall refer whenever necessary.

2. BASICS

2.1. Syntax and semantics. Given a linear order $\mathbb{D} = \langle D, < \rangle$, we call the elements of D *points* (denoted by a, b, \dots) and define an *interval* as an ordered pair $[a, b]$ of points in D , where $a < b$. Abstract intervals will be denoted by I, J, \dots , and so on. Now, as we have mentioned above, there are 13 possible relations, including equality, between any two intervals. From now on, we call these *interval-interval* relations. Besides equality, there are 2 different relations that may hold between any two points (*before* and *after*), called hereafter *point-point* relations, and 5 different relations that may hold between a point and an interval and vice-versa: we call those *interval-point* and *point-interval* relations, respectively, and we use the term *mixed* relations to refer to them indistinctly. Interval-interval relations are exactly Allen’s relations [All83]; point-point relations are the classical relations on a linear order, and mixed relations will be explained below. Traditionally, interval relations are represented by the initial letter of the description of the relation, like m for *meets*. However, when one considers more relations (like point-point and point-interval relations) this notation becomes confusing, and even more so in the presence of more relations, e.g. when one wants to consider interval relations over a *partial order*¹. We introduce the following notation to resolve this issue: an interval $[a, b]$ induces a partition of \mathbb{D} into five regions (see [Lig91]): region 0 which contains all points less than a , region 1 which contains a only, region 2 which contains all the points strictly between a and b , region 3 which contains only b and region 4 which contains the points greater than b . Likewise,

¹This paper is focused on linear orders only; nevertheless, it is our intention to complete this study to include the treatment of partial orders also, and, at this stage, we want to make sure that we will be able to keep the notation consistent.

| | | | | | |
|----------|----------|-------------------------------|---|---|---------|
| $[a, b]$ | 3_{ip} | $c \Leftrightarrow b = c$ | $\begin{array}{c} a \\ \\ \vdots \\ \zeta \\ \\ \vdots \\ \zeta \\ \\ \vdots \end{array}$ | $\begin{array}{c} b \\ \\ \zeta \\ \\ \vdots \\ \zeta \\ \\ \vdots \end{array}$ | ζ |
| $[a, b]$ | 4_{ip} | $c \Leftrightarrow b < c$ | | | |
| $[a, b]$ | 2_{ip} | $c \Leftrightarrow a < c < b$ | | | |
| $[a, b]$ | 1_{ip} | $c \Leftrightarrow a = c$ | | | |
| $[a, b]$ | 0_{ip} | $c \Leftrightarrow c < a$ | ζ | | |

Table 2: Interval-point relations.

a point c induces a partition of \mathbb{D} into 3 pieces: region 0 contains all the points less than c , region 2 contains only c , and region 4 contains all the points greater than c . Interval-interval relations will be denoted by $I\kappa\kappa'_{ii}J$ (where the subscript ii refers to interval-interval relations), where $k, k' \in \{0, 1, 2, 3, 4\}$, and k represents the region of the partition induced by I in which the left endpoint of J falls, while k' is the region of the same partition in which the right endpoint of J falls; for example, $I34_{ii}J$ is exactly Allen’s relation *meets*. Similarly, interval-point relations will be denoted by $I\kappa_{ip}a$ (where the subscript ip stands for interval-point relations), where k represents the position of a with respect to I ; for example, $I4_{ip}a$ is the relation *before*. Analogously, point-point relations will be denoted by the symbol κ_{pp} , and point-interval relations by the symbol $\kappa\kappa'_{pi}$. For point-point relations it is more convenient to use $<$ instead of 4_{pp} , and $>$ instead of 0_{pp} . In Tab. 1 we show six of the interval-interval relations, along with its original nomenclature and symbology, and in Tab. 2 we show the interval-point relations. Finally, we consider one equality per sort, using $=_i$ to denote 13_{ii} (equality between intervals), and $=_p$ to denote 2_{pp} (the equality between points). Now, given any of the mentioned relations r , its inverse, generically denoted by \bar{r} , can be obtained by inverting the roles of the objects in the case of non-mixed relations; for example, the inverse of the relation 22_{ii} (Allen’s relation *contains*) is the relation 04_{ii} (Allen’s relation *during*). On the other hand, mixed relations present a different situation: the inverse of a point-interval relation is an interval-point relation; thus, for example, the inverse of 3_{ip} is 02_{pi} . Finally, notice that some combinations are forbidden: for instance, the relation 22_{pi} makes no sense, as all intervals have a non-zero extension.

Definition 1. We shall denote by: \mathfrak{R} the set of all above described relations; $\mathfrak{I} \subset \mathfrak{R}$ the subset of all 13 interval-interval relations (Allen’s relations) including the relation $=_i$; $\mathfrak{M} \subset \mathfrak{R}$ the subset of all mixed relations; $\mathfrak{P} \subset \mathfrak{R}$ the subset of all point-point relations including the relation $=_p$. Clearly, $\mathfrak{R} = \mathfrak{I} \cup \mathfrak{M} \cup \mathfrak{P}$.

Definition 2. In the following, we denote by:

- (i) **Lin** the class of all linear orders;
- (ii) **Den** the class of all dense linear orders, that is, the class of all linear orders where there exists a point in between any two distinct points;
- (iii) **Dis** the class of all weakly discrete linear orders, that is, the class of all linear orders where each point, other than the least (resp., greatest) point, if there is one, has a direct predecessor (resp., successor) – by a *direct predecessor* of a we of course mean a point b such that $b < a$ and for all points c , if $c < a$ then $c \leq b$, and the notion of a *direct successor* is defined dually;

- (iv) **Unb** the class of all unbounded linear orders, that is, the class of all linear order such that for every point a there exists a point $b > a$ and a point $c < a$.

Definition 3. Given a linear order \mathbb{D} , and given the set $\mathbb{I}(\mathbb{D}) = \{[a, b] \mid a, b \in \mathbb{D}, a < b\}$ of all intervals built on \mathbb{D} :

- a *concrete interval structure of signature S* is a relational structure $\mathcal{F} = \langle \mathbb{I}(\mathbb{D}), r_1, r_2, \dots, r_n \rangle$, where $S = \{r_1, \dots, r_n\} \subseteq \mathfrak{I}$, and
- a *concrete point-interval structure of signature S* is a two-sorted relational structure $\mathcal{F} = \langle \mathbb{D}, \mathbb{I}(\mathbb{D}), r_1, r_2, \dots, r_n \rangle$, where $S = \{r_1, \dots, r_n\} \subseteq \mathfrak{R}$.

Since all relations between intervals, points, and all mixed relations are already implicit in $\mathbb{I}(\mathbb{D})$, we shall often simply write $\langle \mathbb{I}(\mathbb{D}) \rangle$ for a concrete interval structure $\langle \mathbb{I}(\mathbb{D}), r_1, r_2, \dots, r_n \rangle$, and $\langle \mathbb{D}, \mathbb{I}(\mathbb{D}) \rangle$ for a concrete point-interval structure $\langle \mathbb{D}, \mathbb{I}(\mathbb{D}), r_1, r_2, \dots, r_n \rangle$; this is in accordance with the standard usage in much of the literature on interval temporal logics. Moreover, we denote by $FO + S$ the language of first-order logic without equality and relation symbols corresponding to the relations in S . Finally, \mathcal{F} is further said to be *of the class C* ($C \in \{\text{Lin}, \text{Den}, \text{Dis}, \text{Unb}\}$) when \mathbb{D} belongs to the specific class of linear orders C .

2.2. (Un)definability and Truth Preserving Relations. We describe here the most important tools that we use to classify the expressive power of our (sub-)languages.

Definition 4. Let $S \subseteq \mathfrak{R}$, and C a class of linear orders. We say that $FO + S$ *defines* $r \in \mathfrak{R}$ over C , denoted by $FO + S \rightarrow_C r$, if there exists an $FO + S$ -formula $\varphi(x, y)$ such that $\varphi(x, y) \leftrightarrow r(x, y)$ is valid on the class of concrete point-interval structures of signature $(S \cup \{r\})$ based on C .

By $FO + S \rightarrow r$ we denote the fact that $FO + S \rightarrow_{\text{Lin}} r$ (and hence $FO + S \rightarrow_C r$ for every $C \in \{\text{Lin}, \text{Den}, \text{Dis}, \text{Unb}\}$). Obviously, $FO + S \rightarrow r$ for all $r \in S$.

Definition 5. Let $S, S' \subseteq \mathfrak{R}$ and C a class of linear orders. We say that S is:

- *S' -complete over C* (resp., *S' -incomplete over C*) if and only if $FO + S \rightarrow_C r$ for all $r \in S'$ (resp., $FO + S \not\rightarrow_C r$ for some $r \in S'$), and
- *minimally S' -complete over C* (resp., *maximally S' -incomplete over C*) if and only if it is S' -complete (resp., S' -incomplete) over C , and every proper subset (resp., every proper superset) of S is S' -incomplete (resp., S' -complete) over the same class.

The notion of (minimally) r -completeness and (maximally) r -incompleteness over C is immediately deduced from the above one, by taking $S' = \{r\}$ and denoting the latter simply by r . Moreover, one can project the above definitions over some interesting subsets of \mathfrak{R} , such as $\mathfrak{I}, \mathfrak{M}$ or \mathfrak{P} , obtaining relative completeness and incompleteness.

Let $C' \subseteq C$ be two classes of linear orders. Notice that if $FO + S \rightarrow_C r$ then $FO + S \rightarrow_{C'} r$ and, contrapositively, that if $FO + S \not\rightarrow_{C'} r$ then $FO + S \not\rightarrow_C r$. So specifically, if S is S' -complete over C , then it is also S' -complete over C' . Also, if S is S' -incomplete over C' , then it is also S' -incomplete over C . Notice however, that minimality and maximality of complete and incomplete sets does not necessarily transfer between super and subclasses in a similar way. In what follows, in order to prove that $FO + S \not\rightarrow_C r$ for some r and some class C , we shall repeatedly apply the following definition and (rather standard) procedure.

Definition 6. Let $\mathcal{F} = \langle \mathbb{D}, \mathbb{I}(\mathbb{D}), S \rangle$ and $\mathcal{F}' = \langle \mathbb{D}', \mathbb{I}(\mathbb{D}'), S \rangle$ be concrete structures where $S \subseteq \mathfrak{R}$. A binary relation $\zeta \subseteq (\mathbb{D} \cup \mathbb{I}(\mathbb{D})) \times (\mathbb{D}' \cup \mathbb{I}(\mathbb{D}'))$ is called a *surjective S -truth preserving relation* if and only if:

- (i) ζ respects sorts, i.e., $\zeta = \zeta_p \cup \zeta_i$, where $\zeta_p \subseteq \mathbb{D} \times \mathbb{D}'$ and $\zeta_i \subseteq \mathbb{I}(\mathbb{D}) \times \mathbb{I}(\mathbb{D}')$;
- (ii) ζ respects the relations in S , i.e., if $(a, a'), (b, b') \in \zeta_p$ and $(I, I'), (J, J') \in \zeta_i$, then:
 - (a) $r(a, b)$ if and only if $r(a', b')$ for every point-point relation $r \in S$;
 - (b) $r(I, a)$ if and only if $r(I', a')$ for every interval-point relation $r \in S$;
 - (c) $r(I, J)$ if and only if $r(I', J')$ for every interval-interval relation $r \in S$;
- (iii) ζ is total and surjective, i.e.:
 - (a) for every $a \in \mathbb{D}$ (resp., $I \in \mathbb{I}(\mathbb{D})$), there exist $a' \in \mathbb{D}'$ (resp., $I' \in \mathbb{I}(\mathbb{D}')$) such that $(a, a') \in \zeta_p$ (resp., $(I, I') \in \zeta_i$);
 - (b) for every $a' \in \mathbb{D}'$ (resp., $I' \in \mathbb{I}(\mathbb{D}')$), there exist $a \in \mathbb{D}$ (resp., $I \in \mathbb{I}(\mathbb{D})$) such that $(a, a') \in \zeta_p$ (resp., $(I, I') \in \zeta_i$).

If we add to Definition 6 the requirement that ζ should be functional, we obtain nothing but the definition of an isomorphism between two-sorted first-order structures or, equivalently, an isomorphism between single sorted first-order structures with predicates added for ‘point’ and ‘interval’ (see e.g. [Hod93]). As one would expect, surjective S -truth preserving relations preserve the truth of all first-order formulas in signature S . This is stated in Theorem 8, below. The reason why we consider only interval-point relations instead of all mixed relations is that, as we shall explain, we can limit ourselves to work without inverse relations, and point-interval relations are the inverse of interval-point ones.

Definition 7. If ζ is a surjective S -truth preserving relation, we say that ζ *breaks* $r \notin S$ if and only if there are:

- (i) $(a, a'), (b, b') \in \zeta_p$ such that $r(a, b)$ but $\neg r(a', b')$, if r is point-point, or
- (ii) $(a, a') \in \zeta_p$ and $(I, I') \in \zeta_i$ such that $r(I, a)$ but $\neg r(I', a')$, if r is interval-point, or
- (iii) $(I, I'), (J, J') \in \zeta_i$ such that $r(I, J)$ but $\neg r(I', J')$, if r is interval-interval.

The following result is, as already mentioned, a straightforward generalization of the classical result on the preservation of truth under isomorphism between first-order structures, and it is proved by an easy induction on formulas, using clause (ii) of Definition 6 to establish the base case for atomic formulas and clause (iii) for the inductive step for the quantifiers.

Theorem 8. If $\zeta = \zeta_p \cup \zeta_i$ is a surjective S -truth preserving relation between $\mathcal{F} = \langle \mathbb{D}, \mathbb{I}(\mathbb{D}), S \rangle$ and $\mathcal{F}' = \langle \mathbb{D}', \mathbb{I}(\mathbb{D}'), S \rangle$, and $a_1, \dots, a_k \in \mathbb{D}$, $a'_1, \dots, a'_k \in \mathbb{D}$, $I_1, \dots, I_l \in \mathbb{I}(\mathbb{D})$, and $I'_1, \dots, I'_l \in \mathbb{I}(\mathbb{D}')$ are such that $(a_j, a'_j) \in \zeta_p$ for $1 \leq j \leq k$, and $(I_j, I'_j) \in \zeta_i$ for $1 \leq j \leq l$, then for every $FO+S$ formula $\varphi(x_p^1, \dots, x_p^k, y_i^1, \dots, y_i^l)$ with free variables $x_p^1, \dots, x_p^k, y_i^1, \dots, y_i^l$, we have that

$$\mathcal{F} \models \varphi(a_1, \dots, a_k, I_1, \dots, I_l) \text{ if and only if } \mathcal{F}' \models \varphi(a'_1, \dots, a'_k, I'_1, \dots, I'_l).$$

Thus, to show that $FO + S \not\models r$ for a given $r \in \mathfrak{R}$, it is sufficient to find two concrete point-interval structures \mathcal{F} and \mathcal{F}' and a surjective S -truth preserving relation ζ between \mathcal{F} and \mathcal{F}' which breaks r . For the readers' convenience, let us refer to surjective S -truth preserving relations as simply *S -relations*.

Although there are other constructions that could be used to show that relations are not definable in $FO + S$, e.g. elementary embedding or Ehrenfeucht-Fraïssé games, we have found S -relations sufficient for our purposes in this paper.

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|--|--|
| <pre> proc Undef ($r \in \mathfrak{R}^+, def_rules$) for all $S \subseteq \mathfrak{R}^+$ { $S = \text{Closure}(S, def_rules)$; if ($(r \notin S)$ and (S is maximal)) then list S } </pre> | <pre> proc Closure (S, def_rules) while (S changes) { for all $1 \leq i \leq \text{size}(def_rules)$ { if ($def_rules[i]$ applies) { then $S = \text{Apply}(S, def_rule[i])$ } } } return S </pre> |
|--|--|

Figure 1: Pseudo-code to identify maximally r -incomplete sets.

2.3. Strategy. The main objective of this paper is to establish all expressively different subsets of \mathfrak{R} (and, then, of $\mathfrak{I}, \mathfrak{M}$ or \mathfrak{P}) over the mentioned classes of linear orders. To this end, for each $r \in \mathfrak{R}$ we compute all expressively different minimally r -complete and all maximally r -incomplete subsets of \mathfrak{R} , from which we can easily deduce all expressively different minimally r -complete and maximally r -incomplete subsets of $\mathfrak{I}, \mathfrak{M}$ and \mathfrak{P} ; minimally \mathfrak{R} - (resp., \mathfrak{I} -, \mathfrak{M} -, \mathfrak{P} -) complete and maximally incomplete subsets are, then, deduced as a consequence of the above results. The set \mathfrak{R} contains, as we have mentioned, 26 different relations. This means that there are 2^{26} potentially different extensions of first-order logic to be studied. Clearly, unless we design a precise strategy that allows us to reduce the number of results to be proved, the task becomes cumbersome.

As a first simplification principle observe that, since we are working within first-order logic, all inverses of relations are explicitly definable, and hence we only need to assume as primitive a set which contains all relation up to inverses, which implies that point-interval relations can be omitted if we consider all interval-point ones. Accordingly, let \mathfrak{I}^+ be the set of interval-interval relations given in Tab. 1 together with $=_i$, \mathfrak{M}^+ be the set of interval-point relations given in Tab. 2, and let $\mathfrak{P}^+ = \{<, =_p\}$. Lastly let $\mathfrak{R}^+ = \mathfrak{I}^+ \cup \mathfrak{M}^+ \cup \mathfrak{P}^+$.

In order to further reduce the number of results to be presented, consider what follows. The *order dual* of a structure $\mathcal{F} = \langle \mathbb{D}, \mathbb{I}(\mathbb{D}) \rangle$ is the structure $\mathcal{F}^\partial = \langle \mathbb{D}^\partial, \mathbb{I}(\mathbb{D}^\partial) \rangle$ based on the order dual \mathbb{D}^∂ (obtained by reversing the order) of the underlying linear order \mathbb{D} . All classes considered in this paper are closed under taking order duals.

Definition 9. The *reversible relations* are exactly the members of the set $\{0_{ip}, 1_{ip}, 3_{ip}, 4_{ip}, 14_{ii}, 03_{ii}\}$. The relations belonging to the complement $\mathfrak{R}^+ \setminus \{0_{ip}, 1_{ip}, 3_{ip}, 4_{ip}, 14_{ii}, 03_{ii}\}$ are called *symmetric*; if, in addition, $r = 2_{ip}$ or $r = 04_{ii}$, then r is said *self-symmetric*. If $r = 0_{ip}$ (resp., $r = 1_{ip}, r = 14_{ii}$), its *reverse* is $r = 4_{ip}$ (resp., $r = 3_{ip}, r = 03_{ii}$), and the other way around. Finally, the *symmetric* S' of a subset $S \subseteq \mathfrak{R}^+$ is obtained by replacing every reversible relation in S with its reverse. We shall use the notation $S \sim S'$ to indicate that sets S and S' are symmetric.

This definition is motivated by the following easily verifiable facts. Let $r \in \mathfrak{R}^+$, \mathcal{F} be a structure, and x and y be elements of \mathcal{F} of the appropriate sorts for r ; then:

- (i) if r is a reversible relation, with reverse r' , then $\mathcal{F} \models r(x, y)$ if and only if $\mathcal{F}^\partial \models r'(x, y)$;
- (ii) if r is self-symmetric, then $\mathcal{F} \models r(x, y)$ if and only if $\mathcal{F}^\partial \models r(x, y)$;
- (iii) if r is a symmetric, but not self-symmetric, relation, then $\mathcal{F} \models r(x, y)$ if and only if $\mathcal{F}^\partial \models r(y, x)$.

The following crucial lemma capitalizes on these facts.

| | |
|------------------------------------|--|
| $\mathbb{D}, \mathbb{D}', \dots$ | (generic) linearly ordered sets |
| x_p, y_p, \dots | first-order variables for points |
| x_i, y_i, \dots | first-order variables for intervals |
| x, y, \dots | first-order variables of any sort |
| <i>before</i> , ... | relations in text are <i>emphasized</i> |
| $\mathcal{F}, \mathcal{F}', \dots$ | (generic) concrete (point-)interval structures |
| S, S', \dots | (generic) subsets of \mathfrak{R} -relations |
| $\zeta (\zeta_p, \zeta_i)$ | surjective relation (for points, for intervals) |
| $Id_p (Id_i)$ | 'identity' relation over points (intervals) |
| C, C' | (generic) class of linearly ordered sets |
| $FO + S \rightarrow_C r$ | S defines r (w.r.t. the class C) |
| $S \sim S'$ | S and S' are symmetric |
| $a \in \mathbb{D}$ | a is a point of D , where $\mathbb{D} = (D, <)$ |
| <u>S</u> | in the text, a new proof case is <u>underlined</u> |
| r | generic relation |
| $\text{mcs} (\text{mcs}(r))$ | minimally complete set (minimally r -complete set) |
| $\text{MIS} (\text{MIS}(r))$ | maximally incomplete set (maximally r -incomplete set) |

Table 3: Notational conventions used in this paper.

Lemma 10. Let $S, S' \subset \mathfrak{R}^+$ be such that $S \sim S'$. If r is a symmetric relation, then $FO + S \rightarrow r$ if and only if $FO + S' \rightarrow r$. Moreover, if r is a reversible relation with reverse r' , then $FO + S \rightarrow r$ if and only if $FO + S' \rightarrow r'$.

Proof. Let $S, S' \subset \mathfrak{R}^+$ such that $S \sim S'$. For any $FO + S$ formula φ that defines a given relation (and, therefore, with exactly two free variables), let φ' denote the formula obtained from φ by replacing every occurrence of a reversible relation with its reverse, and by swapping the arguments of every symmetric, but not self-symmetric, relation (occurrences of every self-symmetric relation are left unchanged). Induction on formulas then shows that $\mathcal{F} \models \varphi(x, y)$ (after substituting x, y with elements of the appropriate sorts) if and only if $\mathcal{F}^\partial \models \varphi'(x, y)$, for any structure \mathcal{F} . The base case of the induction is taken care of by the three observations preceding this lemma. Now, suppose that a $FO + S$ formula $\varphi(x, y)$ defines a symmetric relation r . We claim that φ' also defines r . Let \mathcal{F} be an arbitrary structure of signature $S \cup \{r\}$. Then $\mathcal{F}^\partial \models \varphi(x, y) \leftrightarrow r(x, y)$, and hence $\mathcal{F} \models \varphi'(x, y) \leftrightarrow r(y, x)$ if r is not self-symmetric, and $\mathcal{F} \models \varphi'(x, y) \leftrightarrow r(x, y)$ otherwise. Next, suppose that the $FO + S$ formula $\varphi(x, y)$ defines a reversible relation r . We claim that φ' defines its reverse r' . Let \mathcal{F} be an arbitrary structure of signature $S \cup \{r\}$. Then $\mathcal{F}^\partial \models \varphi(x, y) \leftrightarrow r(x, y)$, and, hence, $\mathcal{F} \models \varphi'(x, y) \leftrightarrow r'(x, y)$. \square

In conclusion, we can limit our attention to 14 out of 26 relations by disregarding the inverses of relations in \mathfrak{R}^+ , and we do not need to explicitly analyze complete and incomplete sets for 3_{ip} , 4_{ip} , and 03_{ii} as those correspond exactly to symmetric of complete and incomplete sets for 0_{ip} , 1_{ip} , and 14_{ii} , respectively. This means that only 11 relations are to be analyzed (which we can refer to as *explicit* relations).

Even under the mentioned simplifications, there is a huge number of results to be presented and displayed. Let r be anyone of the explicit relations. In order to correctly identifying all minimally r -complete sets ($\text{mcs}(r)$), we need to know all maximally r -incomplete

| $=_p$ | $=_i$ | $<$ | 0_{ip} | 1_{ip} | 2_{ip} |
|--------------|----------------------|--------------------------------|-----------------------------------|-----------------------------------|-----------------------------------|
| $\{<\}$ | $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}, 1_{ip}\}$ | $\{1_{ip}, <\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ |
| $\{1_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{1_{ip}, 2_{ip}\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ |
| $\{3_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ | $\{1_{ip}, 2_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ | $\{2_{ip}, 4_{ip}, <\}$ | $\{1_{ip}, 03_{ii}\}$ |
| | $\{1_{ip}, 2_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ | $\{1_{ip}, 4_{ip}\}$ | $\{2_{ip}, 14_{ii}, <\}$ | $\{3_{ip}, 14_{ii}\}$ |
| | $\{1_{ip}, 3_{ip}\}$ | $\{3_{ip}, 4_{ip}\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{2_{ip}, 03_{ii}, <\}$ | $\{0_{ip}, 3_{ip}\}$ |
| | $\{1_{ip}, 4_{ip}\}$ | $\{1_{ip}, 4_{ip}\}$ | $\{2_{ip}, 4_{ip}\}$ | $\{2_{ip}, 34_{ii}, <\}$ | $\{0_{ip}, 14_{ii}, 24_{ii}, <\}$ |
| | $\{2_{ip}, 3_{ip}\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{2_{ip}, 14_{ii}, <\}$ | $\{3_{ip}, 14_{ii}\}$ | $\{0_{ip}, 14_{ii}, 04_{ii}, <\}$ |
| | $\{2_{ip}, 4_{ip}\}$ | $\{3_{ip}, 14_{ii}\}$ | $\{2_{ip}, 03_{ii}, <\}$ | $\{3_{ip}, 24_{ii}, 03_{ii}\}$ | $\{0_{ip}, 14_{ii}, 44_{ii}, <\}$ |
| | $\{14_{ii}\}$ | $\{3_{ip}, 34_{ii}\}$ | $\{3_{ip}, 14_{ii}\}$ | $\{3_{ip}, 03_{ii}, 04_{ii}\}$ | $\{0_{ip}, 03_{ii}, <\}$ |
| | $\{03_{ii}\}$ | $\{1_{ip}, 03_{ii}\}$ | $\{4_{ip}, 14_{ii}, <\}$ | $\{3_{ip}, 03_{ii}, 44_{ii}\}$ | $\{0_{ip}, 34_{ii}, <\}$ |
| | $\{34_{ii}\}$ | $\{1_{ip}, 34_{ii}\}$ | $\{1_{ip}, 14_{ii}, 24_{ii}\}$ | $\{3_{ip}, 34_{ii}\}$ | $\{1_{ip}, 4_{ip}\}$ |
| | | $\{1_{ip}, 04_{ii}, =_i\}$ | $\{1_{ip}, 14_{ii}, 04_{ii}\}$ | $\{4_{ip}, 14_{ii}, <\}$ | $\{1_{ip}, 14_{ii}, 24_{ii}\}$ |
| | | $\{1_{ip}, 24_{ii}, =_i\}$ | $\{1_{ip}, 14_{ii}, 44_{ii}\}$ | $\{4_{ip}, 24_{ii}, 03_{ii}, <\}$ | $\{1_{ip}, 14_{ii}, 04_{ii}\}$ |
| | | $\{1_{ip}, 44_{ii}, =_i\}$ | $\{1_{ip}, 24_{ii}, =_i\}$ | $\{4_{ip}, 03_{ii}, 04_{ii}, <\}$ | $\{1_{ip}, 14_{ii}, 44_{ii}\}$ |
| | | $\{3_{ip}, 04_{ii}, =_i\}$ | $\{1_{ip}, 04_{ii}, =_i\}$ | $\{4_{ip}, 03_{ii}, 44_{ii}, <\}$ | $\{1_{ip}, 34_{ii}\}$ |
| | | $\{3_{ip}, 24_{ii}, =_i\}$ | $\{1_{ip}, 44_{ii}, =_i\}$ | $\{4_{ip}, 34_{ii}, <\}$ | $\{3_{ip}, 24_{ii}, 03_{ii}\}$ |
| | | $\{3_{ip}, 44_{ii}, =_i\}$ | $\{1_{ip}, 03_{ii}\}$ | | $\{3_{ip}, 03_{ii}, 04_{ii}\}$ |
| | | $\{1_{ip}, 14_{ii}, 24_{ii}\}$ | $\{1_{ip}, 34_{ii}\}$ | | $\{3_{ip}, 03_{ii}, 44_{ii}\}$ |
| | | $\{1_{ip}, 14_{ii}, 04_{ii}\}$ | $\{2_{ip}, 34_{ii}, <\}$ | | $\{3_{ip}, 34_{ii}\}$ |
| | | $\{1_{ip}, 14_{ii}, 44_{ii}\}$ | $\{3_{ip}, 34_{ii}\}$ | | $\{4_{ip}, 14_{ii}, <\}$ |
| | | $\{3_{ip}, 24_{ii}, 03_{ii}\}$ | $\{3_{ip}, 24_{ii}, 03_{ii}\}$ | | $\{4_{ip}, 24_{ii}, 03_{ii}, <\}$ |
| | | $\{3_{ip}, 03_{ii}, 04_{ii}\}$ | $\{3_{ip}, 03_{ii}, 04_{ii}\}$ | | $\{4_{ip}, 03_{ii}, 04_{ii}, <\}$ |
| | | $\{3_{ip}, 03_{ii}, 44_{ii}\}$ | $\{3_{ip}, 03_{ii}, 44_{ii}\}$ | | $\{4_{ip}, 03_{ii}, 44_{ii}, <\}$ |
| | | | $\{4_{ip}, 34_{ii}, <\}$ | | |
| | | | $\{4_{ip}, 24_{ii}, 03_{ii}, <\}$ | | |
| | | | $\{4_{ip}, 03_{ii}, 04_{ii}, <\}$ | | |
| | | | $\{4_{ip}, 03_{ii}, 44_{ii}, <\}$ | | |

Table 4: The spectrum of the $mcs(r)$, for each $r \in \mathfrak{M}^+ \cup \{=_p, =_i, <\}$. - Class: Lin (review).

sets ($MIS(r)$) over the same class, and the other way around. To this end, we proceed in the following way:

- (1) fixed a class of linearly ordered sets and an explicit relation r , we first guess the r -complete subsets of \mathfrak{R}^+ , obtaining a first approximation of the definability rules for r (here, denoted by *def_rules*);
- (2) then, we apply the algorithm in Fig. 1, which uses the set of r -complete subsets of \mathfrak{R}^+ (the parameter *def_rules*) to obtain a first approximation of the maximally r -incomplete sets (the procedure *Closure()* returns the transitive closure of a set of relations, obtained by systematically applying the definition rules contained in *def_rules* - function *Apply()*);
- (3) after that, we prove that every R_1, R_2, \dots, R_k listed as a maximally r -incomplete set is actually r -incomplete, and, if not, we repeat from step 1, using the acquired knowledge to update the set of r -complete subsets of \mathfrak{R}^+ ;
- (4) at this point, the sets $S_1, S_2, \dots, S_{k'}$ listed at step 1 are, actually, all minimally r -complete. To see this, observe that, for each i , S_i is r -complete by definition, and if there was a r -complete set $S \subset S_i$, then for some R_j listed as maximally r -incomplete set we could not prove its r -incompleteness. Therefore, $S_1, S_2, \dots, S_{k'}$

are, in fact, all minimally r -complete, and, as a consequence, R_1, R_2, \dots, R_k are all maximally r -incomplete.

For example, suppose that we are interested in minimal $=_i$ -complete (resp., maximal $=_i$ -incomplete) sets of relations. First, we guess that having the relations $1_{ip}, 2_{ip}$, and 3_{ip} is sufficient to define equality between intervals, because we can easily define the latter by asserting that two intervals are the same if they *start* with the same point, *end* with the same point, and *contain* the same points; so, $S = \{1_{ip}, 2_{ip}, 3_{ip}\}$ is our first guess. Now, by computing all possible subsets, we find that if S were minimal, then $\mathfrak{R}^+ \setminus S$ should be maximally $=_i$ -incomplete. Since we cannot find a proof of the latter, we realize that, in fact, $S' = \{1_{ip}, 3_{ip}\}$ is sufficient to define $=_i$. This back-and-forth procedure continues until it stabilizes; clearly, one cannot find one minimal r -complete subset without finding, at the same time, all maximal r -incomplete subsets, and the other way around. In the following, for each completeness result, its minimality is left unproven; after the corresponding undefinability results are proven, then, by observing that our sets are systematically computed, their respective minimality and maximality will be a consequence.

The most common notational conventions used in the paper are listed in Tab. 3.

3. COMPLETENESS RESULTS IN THE CLASS Den

As it turns out, only a few of the results that appear in Tab. 4 and Tab. 5 are to be refined to obtain all minimally complete sets under the assumption of density. In the following, we show in boldface those complete sets that are stricter than some complete set for a given relation, or completely new, and we prove them explicitly modulo symmetry.

3.1. Definability in \mathfrak{J}^+ and in \mathfrak{M}^+ . We begin our study of definability over the Class Den by considering minimal definability of \mathfrak{J}^+ relations by sets of \mathfrak{J}^+ relations, and then of \mathfrak{M}^+ relations by sets of \mathfrak{M}^+ relations.

Lemma 11. Tab. 6 is correct.

Proof. We begin by proving that $\{24_{ii}\}$ is 34_{ii} -complete. In Part I we proved that on linear domains the relation $34_{ii} \vee 44_{ii}$ is 34_{ii} -complete. It is therefore sufficient to show that $\{24_{ii}\}$ is $(34_{ii} \vee 44_{ii})$ -complete over dense structures. To this end, consider the following definition:

$$x_i 34_{ii} \vee 44_{ii} y_i \leftrightarrow \neg(x_i 24_{ii} y_i) \wedge \exists z_i(x_i 24_{ii} z_i \wedge z_i 24_{ii} y_i). \quad \{24_{ii}\}$$

If $\mathcal{F} \models \varphi([a, b], [c, d])$, there must be some interval z_i *overlapped* by x_i and *overlapping* y_i , which implies $a < c$ and $b < d$ and since x_i cannot *overlap* y_i , we obtain $b \leq c$ as required. Conversely, assume that $[a, b] 34_{ii} \vee 44_{ii} [c, d]$. Then using the density assumption we can take $z_i = [e, f]$, where $a < e < b$ and $c < f < d$, to witness φ . Next, we prove that $\{44_{ii}\}$ is 14_{ii} -complete. Consider the following definition:

$$x_i 14_{ii} y_i \leftrightarrow \forall z_i(z_i 44_{ii} x_i \leftrightarrow z_i 44_{ii} y_i) \wedge \exists z_i(x_i 44_{ii} z_i \wedge \neg(y_i 44_{ii} z_i)). \quad \{44_{ii}\}$$

First suppose that $\mathcal{F} \models \varphi([a, b], [c, d])$. If $a < c$, then we can find a point e with $a < e < c$ and the first conjunct of φ fails, witnessed by the interval $z_i = [a, e]$. Similarly $a > c$

| 34_{ii} | 14_{ii} | 24_{ii} | 04_{ii} | 44_{ii} |
|------------------------|------------------------|--------------------------------|--------------------------------|--------------------------------|
| $\{1_{ip}, 3_{ip}\}$ | $\{0_{ip}, 2_{ip}\}$ | $\{2_{ip}, <\}$ | $\{0_{ip}, 24_{ii}, 44_{ii}\}$ | $\{0_{ip}, 24_{ii}, 04_{ii}\}$ |
| $\{2_{ip}, 14_{ii}\}$ | $\{0_{ip}, 4_{ip}\}$ | $\{0_{ip}, 44_{ii}, 04_{ii}\}$ | $\{1_{ip}, 24_{ii}, 44_{ii}\}$ | $\{1_{ip}, 24_{ii}, 04_{ii}\}$ |
| $\{4_{ip}, 14_{ii}\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{1_{ip}, 44_{ii}, 04_{ii}\}$ | $\{2_{ip}, 24_{ii}, 44_{ii}\}$ | $\{4_{ip}, 24_{ii}, 04_{ii}\}$ |
| $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}, 03_{ii}\}$ | $\{4_{ip}, 44_{ii}, 04_{ii}\}$ | $\{4_{ip}, 24_{ii}, 44_{ii}\}$ | $\{3_{ip}, 24_{ii}, 04_{ii}\}$ |
| $\{2_{ip}, 03_{ii}\}$ | $\{1_{ip}, 2_{ip}\}$ | $\{3_{ip}, 44_{ii}, 04_{ii}\}$ | $\{3_{ip}, 24_{ii}, 44_{ii}\}$ | $\{0_{ip}, 2_{ip}\}$ |
| $\{0_{ip}, 03_{ii}\}$ | $\{1_{ip}, 3_{ip}\}$ | $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ |
| $\{2_{ip}, 4_{ip}\}$ | $\{1_{ip}, 4_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ |
| $\{24_{ii}, 14_{ii}\}$ | $\{1_{ip}, 03_{ii}\}$ | $\{0_{ip}, 4_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ | $\{0_{ip}, 03_{ii}\}$ |
| $\{14_{ii}, 44_{ii}\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{0_{ip}, 03_{ii}\}$ | $\{0_{ip}, 03_{ii}\}$ | $\{1_{ip}, 2_{ip}\}$ |
| $\{14_{ii}, 03_{ii}\}$ | $\{2_{ip}, 4_{ip}\}$ | $\{1_{ip}, 2_{ip}\}$ | $\{1_{ip}, 2_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ |
| $\{14_{ii}, 04_{ii}\}$ | $\{2_{ip}, 03_{ii}\}$ | $\{1_{ip}, 3_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ | $\{1_{ip}, 4_{ip}\}$ |
| $\{24_{ii}, 03_{ii}\}$ | $\{24_{ii}, 03_{ii}\}$ | $\{1_{ip}, 4_{ip}\}$ | $\{1_{ip}, 4_{ip}\}$ | $\{1_{ip}, 03_{ii}\}$ |
| $\{03_{ii}, 04_{ii}\}$ | $\{03_{ii}, 04_{ii}\}$ | $\{1_{ip}, 03_{ii}\}$ | $\{1_{ip}, 03_{ii}\}$ | $\{2_{ip}, 3_{ip}\}$ |
| $\{03_{ii}, 44_{ii}\}$ | $\{03_{ii}, 44_{ii}\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{2_{ip}, 4_{ip}\}$ |
| $\{0_{ip}, 3_{ip}\}$ | $\{34_{ii}\}$ | $\{2_{ip}, 4_{ip}\}$ | $\{2_{ip}, 4_{ip}\}$ | $\{2_{ip}, 14_{ii}\}$ |
| $\{0_{ip}, 4_{ip}\}$ | | $\{2_{ip}, 14_{ii}\}$ | $\{2_{ip}, 14_{ii}\}$ | $\{2_{ip}, 03_{ii}\}$ |
| $\{1_{ip}, 2_{ip}\}$ | | $\{2_{ip}, 03_{ii}\}$ | $\{2_{ip}, 03_{ii}\}$ | $\{3_{ip}, 14_{ii}\}$ |
| $\{1_{ip}, 4_{ip}\}$ | | $\{3_{ip}, 14_{ii}\}$ | $\{2_{ip}, 44_{ii}, <\}$ | $\{4_{ip}, 14_{ii}\}$ |
| $\{1_{ip}, 03_{ii}\}$ | | $\{4_{ip}, 14_{ii}\}$ | $\{3_{ip}, 14_{ii}\}$ | $\{14_{ii}, 24_{ii}\}$ |
| $\{2_{ip}, 3_{ip}\}$ | | $\{14_{ii}, 03_{ii}\}$ | $\{4_{ip}, 14_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ |
| $\{3_{ip}, 14_{ii}\}$ | | $\{14_{ii}, 04_{ii}\}$ | $\{14_{ii}, 24_{ii}\}$ | $\{14_{ii}, 04_{ii}\}$ |
| | | $\{14_{ii}, 44_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ | $\{24_{ii}, 03_{ii}\}$ |
| | | $\{03_{ii}, 44_{ii}\}$ | $\{14_{ii}, 44_{ii}\}$ | $\{03_{ii}, 04_{ii}\}$ |
| | | $\{04_{ii}, 03_{ii}\}$ | $\{24_{ii}, 03_{ii}\}$ | $\{34_{ii}\}$ |
| | | $\{34_{ii}\}$ | $\{03_{ii}, 44_{ii}\}$ | |
| | | | $\{34_{ii}\}$ | |

Table 5: The spectrum of the $mcs(r)$, for each $r \in \mathcal{J}^+ \setminus \{=i\}$. - Class: Lin (review).

| 34_{ii} | 14_{ii} | 24_{ii} | 04_{ii} | 44_{ii} | $=i$ |
|------------------------|------------------------|------------------------|------------------------|------------------------|---------------|
| $\{14_{ii}, 03_{ii}\}$ | $\{34_{ii}\}$ | $\{34_{ii}\}$ | $\{34_{ii}\}$ | $\{34_{ii}\}$ | $\{34_{ii}\}$ |
| $\{14_{ii}, 04_{ii}\}$ | $\{24_{ii}\}$ | $\{44_{ii}\}$ | $\{24_{ii}\}$ | $\{24_{ii}\}$ | $\{14_{ii}\}$ |
| $\{24_{ii}\}$ | $\{03_{ii}, 04_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ | $\{44_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ | $\{03_{ii}\}$ |
| $\{03_{ii}, 04_{ii}\}$ | $\{44_{ii}\}$ | $\{14_{ii}, 04_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ | $\{14_{ii}, 04_{ii}\}$ | |
| $\{44_{ii}\}$ | | $\{03_{ii}, 04_{ii}\}$ | | $\{03_{ii}, 04_{ii}\}$ | |

Table 6: The spectrum of the $mcs(r)$, for each $r \in \mathcal{J}^+$. - Class: Den.
(Lemma 11)

also gives a contradiction, so we obtain $a = c$. The second conjunct of φ gives an interval $z_i = [e, f]$ such that $b < e$ and $e \leq d$, and so we have $b < d$. On the other hand, if we assume that $a = c < b < d$, since the underlying linear order is dense, there exists e such that $b < e < d$; by taking $z_i = [e, d]$, we witness the second conjunct of φ , while the first conjunct is satisfied trivially. The 34_{ii} -completeness of $\{44_{ii}\}$ now follows from the just-proved 14_{ii} -completeness of $\{44_{ii}\}$ and the 34_{ii} -completeness of $\{14_{ii}, 44_{ii}\}$ over all linear orders (see Table 5). Having established this, all the remaining new completeness results listed in Table 6 follow via the completeness of $\{34_{ii}\}$ with respect to every \mathcal{J}^+ -relation. \square

| 0_{ip} | 1_{ip} | 2_{ip} |
|----------------------|----------------------|----------------------|
| $\{1_{ip}, 2_{ip}\}$ | $\{0_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ |
| $\{1_{ip}, 3_{ip}\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ |
| $\{2_{ip}, 4_{ip}\}$ | $\{2_{ip}, 4_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ |
| $\{2_{ip}, 3_{ip}\}$ | | $\{1_{ip}, 4_{ip}\}$ |

Table 7: The spectrum of the $mcs(r)$, for each $r \in \mathfrak{M}^+$. - Class: Den.
(Lemma 14)

The next two lemmas involve both interval-interval and interval-point relations, and therefore contrast with the other results of the present subsection that all concern definability purely among interval-interval relations and purely among interval-point relations. However, anticipating these particular results here, which do not refer to any table in particular, will be very useful in order to simplify some definitions and chains of deduction further on.

Lemma 12. The set $\{0_{ip}\}$ is $<$ -complete over dense linear orders.

Proof. Consider the following definition:

$$x_p < y_p \leftrightarrow (\neg \exists x_i(x_i 0_{ip} y_p) \wedge \exists y_i(y_i 0_{ip} x_p)) \vee \exists x_i(x_i 0_{ip} x_p \wedge \neg(x_i 0_{ip} y_p)) \quad \{0_{ip}\}$$

Suppose that $\mathcal{F} \models \varphi(x_p, y_p)$. If the first disjunct of φ holds then y_p is the greatest point of the model, while x_p is not, which gives $x_p < y_p$. Assume the second disjunct of φ holds. Then the interval $x_i = [a, b]$ is such that $x_p = c < a$, but $y_p \geq a$, that is $x_p < y_p$. Conversely, suppose that $x_p = a, y_p = b$, and that $a < b$. As the structure is dense, if b is the last point of the model, we have no interval *starting* at b but infinitely many intervals x_i such that $x_i 0_{ip} a$, and if b is not the last point of the model, then every x_i *starting* at b satisfies the second part of the definition. \square

Lemma 13. The sets $\{3_{ip}, 04_{ii}\}$ and $\{4_{ip}, 04_{ii}\}$ are 34_{ii} -complete over dense linear orders..

Proof. We will show that both sets $\{3_{ip}, 04_{ii}\}$ and $\{4_{ip}, 04_{ii}\}$ define 03_{ii} in the dense case; then, since $\{03_{ii}, 04_{ii}\}$ defines 34_{ii} in the general case, we have the result. Indeed, consider the following definitions:

$$x_i 03_{ii} y_i \leftrightarrow \begin{cases} \exists z_p(x_i 3_{ip} z_p \wedge y_i 3_{ip} z_p) \wedge \exists z_i(z_i 04_{ii} y_i \wedge \neg(z_i 04_{ii} x_i)). & \{3_{ip}, 04_{ii}\} \\ \forall z_p(x_i 4_{ip} z_p \leftrightarrow y_i 4_{ip} z_p) \wedge \exists z_i(z_i 04_{ii} y_i \wedge \neg(z_i 04_{ii} x_i)). & \{4_{ip}, 04_{ii}\} \end{cases}$$

Let us consider $\{3_{ip}, 04_{ii}\}$, first. If $[a, b] 03_{ii} [c, d]$ then $c < a$ and $b = d$, so, taking z_p equal to b satisfies the first conjunct. By density there is a point e such that $c < e < a$, and therefore taking z_i equal to $[e, a]$ satisfies the second conjunct. Conversely, suppose that $\varphi([a, b], [c, d])$. The first conjunct of φ ensures that $b = d$. If it were the case that $a \leq c$, every interval contained in $[a, b]$ would be contained in $[c, d]$, violating the second conjunct, so it can only be the case that $c < a$. We therefore have that $[a, b] 03_{ii} [c, d]$, as desired. The argument for $\{4_{ip}, 04_{ii}\}$ is similar, using the fact that $\forall z_p(x_i 4_{ip} z_p \leftrightarrow y_i 4_{ip} z_p)$ is true if and only if the intervals assigned x_i and y_i have the same end point. \square

| $=_p$ | $=_i$ | $<$ |
|--------------|----------------------|-----------------------|
| $\{<\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{0_{ip}\}$ |
| $\{0_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ | $\{1_{ip}, 2_{ip}\}$ |
| $\{1_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ |
| $\{3_{ip}\}$ | $\{1_{ip}, 4_{ip}\}$ | $\{2_{ip}, 3_{ip}\}$ |
| $\{4_{ip}\}$ | $\{2_{ip}\}$ | $\{3_{ip}, 14_{ii}\}$ |
| | $\{14_{ii}\}$ | $\{3_{ip}, 34_{ii}\}$ |
| | $\{03_{ii}\}$ | $\{1_{ip}, 03_{ii}\}$ |
| | $\{34_{ii}\}$ | $\{1_{ip}, 34_{ii}\}$ |
| | $\{24_{ii}\}$ | $\{1_{ip}, 04_{ii}\}$ |
| | $\{04_{ii}\}$ | $\{1_{ip}, 24_{ii}\}$ |
| | $\{44_{ii}\}$ | $\{1_{ip}, 44_{ii}\}$ |
| | | $\{3_{ip}, 04_{ii}\}$ |
| | | $\{3_{ip}, 24_{ii}\}$ |
| | | $\{3_{ip}, 44_{ii}\}$ |
| | | $\{4_{ip}\}$ |

Table 8: The spectrum of the $mcs(r)$, for each $r \in \{=_p, =_i, <\}$. - Class: Den.

Lemma 14. Tab. 7 is correct.

Proof. First, we show that $\{0_{ip}\}$ is 1_{ip} -complete. From Lemma 12 we know that $\{0_{ip}\}$ defines $<$ in the dense case, so we may consider the following definition:

$$x_i 1_{ip} y_p \leftrightarrow \forall z_p (x_i 0_{ip} z_p \leftrightarrow z_p < y_p). \quad \{0_{ip}\}$$

If $[a, b] 1_{ip} c$, then $a = c$, and hence $\varphi([a, b], c)$. Conversely, suppose that $\varphi([a, b], c)$. So, if $c < a$, then $[a, b] 0_{ip} c$ by definition of 0_{ip} and so, by $\varphi([a, b], c)$ we obtain the contradiction $c < c$. Now, if $a < c$ then, by density, there is a point d such that $a < d < c$. Hence $\neg[a, b] 0_{ip} d$ by the definition of 0_{ip} , but then $d \not< c$ by $\varphi([a, b], c)$ — a contradiction. The only possible case is therefore that $a = c$, i.e., that $[a, b] 1_{ip} c$. To see that $\{2_{ip}, 4_{ip}\}$ is 1_{ip} -complete, it suffices to note that $\{2_{ip}, 4_{ip}\}$ defines 0_{ip} in the general case, and then appeal to the just proven fact that $\{0_{ip}\}$ defines 1_{ip} . \square

3.2. Definability for the Relations $=_p, =_i, <$ in Den.

Lemma 15. Tab. 8 is correct.

Proof. We begin by showing that $\{0_{ip}\}$ is $=_p$ -complete. Consider the following definition:

$$x_p =_p y_p \leftrightarrow \forall x_i (x_i 0_{ip} x_p \leftrightarrow x_i 0_{ip} y_p) \quad \{0_{ip}\}$$

Suppose, first, that $x_p =_p y_p$; we have that, either there exists no interval in the model, in which case the right-hand side of the definition is vacuously true, or there are infinitely many intervals (as the underlying linear order is dense), and every interval x_i clearly has both x_p and y_p in relation 0_{ip} with it or none of them, again satisfying the right-hand side. If, on the other hand, $\varphi(x_p, y_p)$ is satisfied, then, if there is no interval, then x_p and y_p coincide, and if there are infinitely many intervals, the only way to guarantee that the

intervals that see both x_p and y_p via 0_{ip} are exactly the same is to assign the same point to x_p and y_p , as we wanted to prove. The $=_p$ -completeness of $\{4_{ip}\}$ follows by symmetry (see Lemma 10). Now, consider $=_i$ -completeness, and the following definitions:

$$x_i =_i y_i \leftrightarrow \begin{cases} \forall x_p(x_i 2_{ip} x_p \leftrightarrow y_i 2_{ip} x_p) & \{2_{ip}\} \\ \forall z_i(x_i 24_{ii} z_i \leftrightarrow y_i 24_{ii} z_i) \wedge \forall z_i(z_i 24_{ii} x_i \leftrightarrow z_i 24_{ii} y_i) & \{24_{ii}\} \\ \forall z_i(x_i 04_{ii} z_i \leftrightarrow y_i 04_{ii} z_i) \wedge \forall z_i(z_i 04_{ii} x_i \leftrightarrow z_i 04_{ii} y_i) & \{04_{ii}\} \\ \forall z_i(x_i 44_{ii} z_i \leftrightarrow y_i 44_{ii} z_i) \wedge \forall z_i(z_i 44_{ii} x_i \leftrightarrow z_i 44_{ii} y_i). & \{44_{ii}\} \end{cases}$$

As for $\{2_{ip}\}$, observe that two intervals over a dense linear order are equal if and only if they have the same internal points. Consider, now, the set $\{24_{ii}\}$. It is clear that if $x_i =_i y_i$ then $\varphi(x_i, y_i)$ holds. For the converse, suppose that $[a, b] \neq [c, d]$. Then either $[a, b]$ has an internal point which is not an internal point of $[c, d]$ or $[c, d]$ has an internal point which is not an internal point of $[a, b]$. In both cases either the first or the second conjunct of φ does not hold. The remaining cases are treated with similar arguments. Let us now focus on the relation $<$ and the corresponding definitions:

$$x_p < y_p \leftrightarrow \exists x_i(x_i 0_{ip} x_p \wedge \neg(x_i 0_{ip} y_p)) \quad \{0_{ip}\}$$

Consider the set $\{0_{ip}\}$. Suppose that $\mathcal{F} \models \varphi(x_p, y_p)$. Then, there must be an interval that starts after x_p but not after y_p ; if $x_p \geq y_p$, then it is impossible to find such an interval, so $x_p < y_p$. Conversely, suppose that $x_p = a, y_p = b$, and that $a < b$. As the structure is dense, there exists a point c such that $a < c < b$, and therefore the interval $[c, b]$ can be used as a witness of x_i . All other new definitions follow directly from the results for the class Lin combined with the density hypothesis. In particular, to see that $\{1_{ip}, 04_{ii}\}$ is $<$ -complete, note that $\{04_{ii}\}$ defines $=_i$ in the dense case and that $\{1_{ip}, 04_{ii}, =_i\}$ defines $<$ in the general one. The case for $\{1_{ip}, 24_{ii}\}$ follows since $\{24_{ii}\}$ defines $=_i$ in the dense case and since $\{1_{ip}, 24_{ii}, =_i\}$ defines $<$. The case $\{1_{ip}, 44_{ii}\}$ follows since $\{44_{ii}\}$ defines $=_i$ in the dense case and $\{1_{ip}, 44_{ii}, =_i\}$ defines $<$. The case $\{3_{ip}, 04_{ii}\}$ follows since, in the dense case, we have that $\{04_{ii}\}$ defines $=_i$ and that $\{3_{ip}, 04_{ii}, =_i\}$ defines $<$. The case $\{3_{ip}, 24_{ii}\}$ follows since, under the density hypothesis, we have that $\{24_{ii}\}$ defines $=_i$ and $\{3_{ip}, 24_{ii}, =_i\}$ defines $<$. The case $\{3_{ip}, 44_{ii}\}$ follows since $\{44_{ii}\}$ defines $=_i$ in the dense case and since we already know that $\{3_{ip}, 44_{ii}, =_i\}$ defines $<$. Lastly, the fact that $\{4_{ip}\}$ is $<$ -complete follows from $\{0_{ip}\}$ being able to define $<$ and from Lemma 10. \square

3.3. Definability for \mathfrak{M}^+ -Relations in Den. To continue studying how the density hypothesis influences the ability of the sub-languages to express our relations, we focus now on mixed relations.

Lemma 16. Tab. 9 is correct.

| 0_{ip} | 1_{ip} | 2_{ip} |
|-----------------------|------------------------|-----------------------|
| $\{1_{ip}, <\}$ | $\{0_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ |
| $\{2_{ip}, <\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ |
| $\{1_{ip}, 2_{ip}\}$ | $\{2_{ip}, 4_{ip}\}$ | $\{0_{ip}, 24_{ii}\}$ |
| $\{1_{ip}, 3_{ip}\}$ | $\{2_{ip}, <\}$ | $\{0_{ip}, 44_{ii}\}$ |
| $\{1_{ip}, 4_{ip}\}$ | $\{3_{ip}, 14_{ii}\}$ | $\{0_{ip}, 03_{ii}\}$ |
| $\{2_{ip}, 3_{ip}\}$ | $\{3_{ip}, 24_{ii}\}$ | $\{0_{ip}, 34_{ii}\}$ |
| $\{2_{ip}, 4_{ip}\}$ | $\{3_{ip}, 04_{ii}\}$ | $\{0_{ip}, 04_{ii}\}$ |
| $\{3_{ip}, 14_{ii}\}$ | $\{3_{ip}, 44_{ii}\}$ | $\{1_{ip}, 3_{ip}\}$ |
| $\{4_{ip}, 14_{ii}\}$ | $\{3_{ip}, 34_{ii}\}$ | $\{1_{ip}, 4_{ip}\}$ |
| $\{1_{ip}, 24_{ii}\}$ | $\{4_{ip}, 14_{ii}\}$ | $\{1_{ip}, 24_{ii}\}$ |
| $\{1_{ip}, 04_{ii}\}$ | $\{4_{ip}, 24_{ii},\}$ | $\{1_{ip}, 44_{ii}\}$ |
| $\{1_{ip}, 44_{ii}\}$ | $\{4_{ip}, 04_{ii}\}$ | $\{1_{ip}, 03_{ii}\}$ |
| $\{1_{ip}, 03_{ii}\}$ | $\{4_{ip}, 44_{ii}\}$ | $\{1_{ip}, 34_{ii}\}$ |
| $\{1_{ip}, 34_{ii}\}$ | $\{4_{ip}, 34_{ii}\}$ | $\{1_{ip}, 04_{ii}\}$ |
| $\{3_{ip}, 34_{ii}\}$ | | $\{3_{ip}, 14_{ii}\}$ |
| $\{3_{ip}, 24_{ii}\}$ | | $\{3_{ip}, 24_{ii}\}$ |
| $\{3_{ip}, 04_{ii}\}$ | | $\{3_{ip}, 44_{ii}\}$ |
| $\{3_{ip}, 44_{ii}\}$ | | $\{3_{ip}, 34_{ii}\}$ |
| $\{4_{ip}, 34_{ii}\}$ | | $\{3_{ip}, 04_{ii}\}$ |
| $\{4_{ip}, 24_{ii}\}$ | | $\{4_{ip}, 14_{ii}\}$ |
| $\{4_{ip}, 04_{ii}\}$ | | $\{4_{ip}, 24_{ii}\}$ |
| $\{4_{ip}, 44_{ii}\}$ | | $\{4_{ip}, 44_{ii}\}$ |
| | | $\{4_{ip}, 34_{ii}\}$ |
| | | $\{4_{ip}, 04_{ii}\}$ |

Table 9: The spectrum of the $mcs(r)$, for each $r \in \mathfrak{M}^+$. - Class: Den.

Proof. As for the $\{0_{ip}\}$ -completeness, only one new definition is necessary:

$$x_i 0_{ip} y_p \leftrightarrow \exists k_p (\forall z_p (x_i 2_{ip} z_p \rightarrow k_p < z_p) \wedge y_p < k_p) \quad \{2_{ip}, <\}$$

Let us prove that the set $\{2_{ip}, <\}$ is 0_{ip} -complete. If $\mathcal{F} \models \varphi(x_i, y_p)$, then there exists a point k_p smaller than every point *contained* in x_i (there are infinitely many such points because we are in a dense structure); and y_p is smaller than k_p , so it must be *before* x_i . If, on the other hand $x_i = [a, b]$ and $y_p = c < a$, we take $k_p = a$ to satisfy all requirements. The 0_{ip} -completeness of the remaining new sets is now a consequence of the 0_{ip} -completeness and previously established results. We will give the details of the deductions chains involved. For $\{4_{ip}, 14_{ii}\}$, note that $\{4_{ip}\} \rightarrow_{\text{Den}} <$ (Table 8) and that $\{4_{ip}, 14_{ii}, <\} \rightarrow_{\text{Lin}} 0_{ip}$ (Table 4). For the cases $\{1_{ip}, 24_{ii}\}$, $\{1_{ip}, 04_{ii}\}$ and $\{1_{ip}, 44_{ii}\}$, we note that $\{24_{ii}\}$, $\{04_{ii}\}$ and $\{44_{ii}\}$ all define $=_i$ over Den (Table 8), and then that $\{1_{ip}, 24_{ii}, =_i\}$, $\{1_{ip}, 04_{ii}, =_i\}$ and $\{1_{ip}, 44_{ii}, =_i\}$ all define 0_{ip} over Lin (Table 4). The cases for $\{3_{ip}, 24_{ii}\}$, $\{3_{ip}, 04_{ii}\}$ and $\{3_{ip}, 44_{ii}\}$ follow similarly, using that $\{3_{ip}, 24_{ii}, =_i\}$, $\{3_{ip}, 04_{ii}, =_i\}$ and $\{3_{ip}, 44_{ii}, =_i\}$ all define 0_{ip} over Lin (Table 4). Considering the 0_{ip} -completeness of $\{4_{ip}, 34_{ii}\}$, we again use the fact that $\{4_{ip}\} \rightarrow_{\text{Den}} <$ (Table 8) and that $\{4_{ip}, 34_{ii}, <\} \rightarrow_{\text{Lin}} 0_{ip}$ (Table 4). In the case of $\{4_{ip}, 24_{ii}\}$ the chain of deductions goes as follows: $\{4_{ip}\} \rightarrow_{\text{Den}} <$ (Table 8), $\{24_{ii}\} \rightarrow_{\text{Den}} 14_{ii}$ (Table 6) and hence, by Lemma 10, $\{24_{ii}\} \rightarrow_{\text{Den}} 03_{ii}$, and lastly $\{4_{ip}, 24_{ii}, 03_{ii}, <\} \rightarrow_{\text{Lin}} 0_{ip}$ (Table 4). In the case of

$\{4_{ip}, 04_{ii}\}$, recall that $\{4_{ip}, 04_{ii}\} \rightarrow_{\text{Den}} 34_{ii}$ (Lemma 13), and then appeal to the just proved 0_{ip} -completeness of $\{4_{ip}, 34_{ii}\}$. In the case of $\{4_{ip}, 44_{ii}\}$ we use the fact that $\{44_{ii}\} \rightarrow_{\text{Den}} 34_{ii}$ and again appeal to the case for $\{4_{ip}, 34_{ii}\}$. As for proving that $\{0_{ip}\}$ is 1_{ip} -complete, observe that by the previous lemma this set defines $<$, and consider the following simple definition:

$$x_i 1_{ip} y_p \leftrightarrow \forall z_p (x_i 0_{ip} z_p \leftrightarrow z_p < y_p) \quad \{0_{ip}\}$$

The left to right direction of the above definition is immediate. For the right to left direction, suppose, on the contrary, that the formula holds and it is not the case that $x_i 1_{ip} y_p$. Then, assume that $x_i = [a, b]$: if $y_p < a$, then, since the underlying linear order is dense, we could find c such that $y_p < c < a$, and contradict the right-side by instantiating z_p with c , and if $a < y_p$, then since the underlying linear order is dense, we could find c such that $a < c < y_p$, and, again, contradict the right-side by instantiating z_p with c . Now, the remaining new definitions for the relation 1_{ip} all follow straightforwardly by noting that these sets all define 0_{ip} . Finally, we prove that $\{1_{ip}, 04_{ii}\}$ is 2_{ip} -complete by means of the following definability equation:

$$x_i 2_{ip} y_p \leftrightarrow \exists z_i (z_i 1_{ip} y_p \wedge z_i 04_{ii} x_i). \quad \{1_{ip}, 04_{ii}\}$$

The 2_{ip} -completeness of $\{1_{ip}, 04_{ii}\}$ is extremely simple. Assuming $\mathcal{F} \models \varphi([a, b], c)$, there must be an interval $z_i = [c, d]$ contained in $[a, b]$, which implies $a < c < b$, as we wanted. Conversely, if $a < c < b$, since the structure is dense, we can find d such that $a < c < d < b$, and $z_i = [c, d]$ is a witness for φ . For the 2_{ip} -completeness of $\{0_{ip}, 24_{ii}\}$, $\{0_{ip}, 44_{ii}\}$ and $\{0_{ip}, 34_{ii}\}$ we note that $\{0_{ip}\} \rightarrow_{\text{Den}} 1_{ip}$, that $\{24_{ii}\}$, $\{44_{ii}\}$ and $\{34_{ii}\}$ each define 04_{ii} over Den, and the just proved 2_{ip} -completeness of $\{1_{ip}, 04_{ii}\}$. For $\{0_{ip}, 03_{ii}\}$ we note that $\{0_{ip}\} \rightarrow_{\text{Den}} <$ and that $\{0_{ip}, 03_{ii}, <\} \rightarrow_{\text{Lin}} 2_{ip}$. For $\{1_{ip}, 24_{ii}\}$ and $\{1_{ip}, 44_{ii}\}$ we again use the fact that $\{24_{ii}\}$ and $\{44_{ii}\}$ each define 04_{ii} over Den and that $\{1_{ip}, 04_{ii}\} \rightarrow_{\text{Den}} 02_{ii}$ as proved above. For $\{3_{ip}, 24_{ii}\}$ and $\{3_{ip}, 44_{ii}\}$ we use the facts that $\{24_{ii}\}$ and $\{44_{ii}\}$ each defines 14_{ii} over Den and that $\{3_{ip}, 14_{ii}\} \rightarrow_{\text{Lin}} 2_{ip}$. In the case of $\{3_{ip}, 04_{ii}\}$ we use the facts that $\{3_{ip}, 04_{ii}\} \rightarrow_{\text{Den}} 34_{ii}$ (Lemma 13) and that $\{3_{ip}, 34_{ii}\} \rightarrow_{\text{Lin}} 2_{ip}$. For $\{4_{ip}, 34_{ii}\}$ we use the previously proven facts that $\{4_{ip}, 34_{ii}\} \rightarrow_{\text{Den}} 1_{ip}$ and $\{1_{ip}, 34_{ii}\} \rightarrow_{\text{Lin}} 2_{ip}$. From this the cases for $\{4_{ip}, 24_{ii}\}$, $\{4_{ip}, 44_{ii}\}$ and $\{4_{ip}, 04_{ii}\}$ now respectively follow since $\{24_{ii}\} \rightarrow_{\text{Den}} 34_{ii}$, $\{44_{ii}\} \rightarrow_{\text{Den}} 34_{ii}$ and $\{4_{ip}, 04_{ii}\} \rightarrow_{\text{Den}} 34_{ii}$ (Lemma 13). Lastly the 2_{ip} -completeness of $\{4_{ip}, 14_{ii}\}$ follows by noting that $\{4_{ip}\} \rightarrow_{\text{Den}} <$ and that $\{4_{ip}, 14_{ii}, <\} \rightarrow_{\text{Lin}} 2_{ip}$. \square

3.4. Definability for Relations in $\mathcal{J}^+ \setminus \{=\}$ in Den. To conclude this analysis of dense linear orders, we prove that Tab. 10 is correct.

Lemma 17. Tab. 10 is correct.

Proof. All new definitions involving only interval-interval relations have already been treated in Lemma 11. All new 14_{ii} -completeness results follow from previously proven facts, and we will give the deduction chains. For $\{2_{ip}, <\}$, we have that $\{2_{ip}, <\}$ defines 0_{ip} in the dense case and that $\{0_{ip}, 2_{ip}\}$ defines 14_{ii} . The cases for $\{3_{ip}, 04_{ii}\}$ and $\{4_{ip}, 04_{ii}\}$ follow

| 14_{ii} | 34_{ii} | 24_{ii} | 04_{ii} | 44_{ii} |
|------------------------|------------------------|------------------------|------------------------|------------------------|
| $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{0_{ip}, 2_{ip}\}$ |
| $\{0_{ip}, 3_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ |
| $\{0_{ip}, 4_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ | $\{0_{ip}, 03_{ii}\}$ | $\{0_{ip}, 4_{ip}\}$ |
| $\{0_{ip}, 03_{ii}\}$ | $\{0_{ip}, 03_{ii}\}$ | $\{0_{ip}, 03_{ii}\}$ | $\{1_{ip}, 3_{ip}\}$ | $\{0_{ip}, 03_{ii}\}$ |
| $\{0_{ip}, 04_{ii}\}$ | $\{0_{ip}, 04_{ii}\}$ | $\{0_{ip}, 04_{ii}\}$ | $\{1_{ip}, 4_{ip}\}$ | $\{0_{ip}, 04_{ii}\}$ |
| $\{1_{ip}, 2_{ip}\}$ | $\{1_{ip}, 2_{ip}\}$ | $\{1_{ip}, 2_{ip}\}$ | $\{1_{ip}, 03_{ii}\}$ | $\{1_{ip}, 2_{ip}\}$ |
| $\{1_{ip}, 3_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ | $\{2_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ |
| $\{1_{ip}, 4_{ip}\}$ | $\{1_{ip}, 4_{ip}\}$ | $\{1_{ip}, 4_{ip}\}$ | $\{3_{ip}, 14_{ii}\}$ | $\{1_{ip}, 4_{ip}\}$ |
| $\{1_{ip}, 03_{ii}\}$ | $\{1_{ip}, 03_{ii}\}$ | $\{1_{ip}, 03_{ii}\}$ | $\{4_{ip}, 14_{ii}\}$ | $\{1_{ip}, 03_{ii}\}$ |
| $\{1_{ip}, 04_{ii}\}$ | $\{1_{ip}, 04_{ii}\}$ | $\{1_{ip}, 04_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ | $\{1_{ip}, 04_{ii}\}$ |
| $\{2_{ip}, 3_{ip}\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{24_{ii}\}$ | $\{2_{ip}, 3_{ip}\}$ |
| $\{2_{ip}, 4_{ip}\}$ | $\{2_{ip}, 4_{ip}\}$ | $\{2_{ip}, 4_{ip}\}$ | $\{34_{ii}\}$ | $\{2_{ip}, 4_{ip}\}$ |
| $\{2_{ip}, 03_{ii}\}$ | $\{2_{ip}, 14_{ii}\}$ | $\{2_{ip}, 14_{ii}\}$ | $\{44_{ii}\}$ | $\{2_{ip}, 14_{ii}\}$ |
| $\{2_{ip}, <\}$ | $\{2_{ip}, 03_{ii}\}$ | $\{2_{ip}, 03_{ii}\}$ | | $\{2_{ip}, 03_{ii}\}$ |
| $\{3_{ip}, 04_{ii}\}$ | $\{2_{ip}, <\}$ | $\{2_{ip}, <\}$ | | $\{2_{ip}, <\}$ |
| $\{4_{ip}, 04_{ii}\}$ | $\{3_{ip}, 14_{ii}\}$ | $\{3_{ip}, 14_{ii}\}$ | | $\{3_{ip}, 14_{ii}\}$ |
| $\{24_{ii}\}$ | $\{3_{ip}, 04_{ii}\}$ | $\{3_{ip}, 04_{ii}\}$ | | $\{3_{ip}, 04_{ii}\}$ |
| $\{03_{ii}, 04_{ii}\}$ | $\{4_{ip}, 14_{ii}\}$ | $\{4_{ip}, 14_{ii}\}$ | | $\{4_{ip}, 14_{ii}\}$ |
| $\{34_{ii}\}$ | $\{4_{ip}, 04_{ii}\}$ | $\{4_{ip}, 04_{ii}\}$ | | $\{4_{ip}, 04_{ii}\}$ |
| $\{44_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ | | $\{14_{ii}, 03_{ii}\}$ |
| | $\{14_{ii}, 04_{ii}\}$ | $\{14_{ii}, 04_{ii}\}$ | | $\{14_{ii}, 04_{ii}\}$ |
| | $\{24_{ii}\}$ | $\{03_{ii}, 04_{ii}\}$ | | $\{24_{ii}\}$ |
| | $\{03_{ii}, 04_{ii}\}$ | $\{34_{ii}\}$ | | $\{03_{ii}, 04_{ii}\}$ |
| | $\{44_{ii}\}$ | $\{44_{ii}\}$ | | $\{34_{ii}\}$ |

Table 10: The spectrum of the $mcs(r)$, for each $r \in \mathcal{J}^+ \setminus \{=i\}$. - Class: Den.

from Lemma 13 and the fact that $\{34_{ii}\}$ define 14_{ii} . The new 34_{ii} -completeness results are based on the following arguments. For $\{0_{ip}, 04_{ii}\}$, we have that $\{0_{ip}, 04_{ii}\}$ defines 2_{ip} in the dense case and that $\{0_{ip}, 2_{ip}\}$ defines 34_{ii} . For $\{1_{ip}, 04_{ii}\}$, we know that $\{1_{ip}, 04_{ii}\}$ defines 2_{ip} in the dense case and that $\{1_{ip}, 2_{ip}\}$ defines 34_{ii} . For $\{2_{ip}, <\}$, we have that $\{2_{ip}, <\}$ defines 0_{ip} in the dense case and that $\{0_{ip}, 2_{ip}\}$ is 34_{ii} -complete. The cases for $\{3_{ip}, 04_{ii}\}$ and $\{4_{ip}, 04_{ii}\}$ were proven in Lemma 13. The 24_{ii} -completeness of the sets $\{0_{ip}, 04_{ii}\}$, $\{3_{ip}, 04_{ii}\}$ and $\{4_{ip}, 04_{ii}\}$ follow from their already proven 34_{ii} -completeness over Den and the 24_{ii} -completeness of 34_{ii} over Lin. For the 04_{ii} -completeness of $\{2_{ip}\}$, assuming density, recall that $\{2_{ip}\}$ is $=i$ -complete in the dense case, and consider the following definition:

$$\begin{aligned}
x_i 04_{ii} y_i &\leftrightarrow \forall z_p (x_i 2_{ip} z_p \rightarrow y_i 2_{ip} z_p) \wedge \exists z_p (y_i 2_{ip} z_p \wedge \neg(x_i 2_{ip} z_p)) && \{2_{ip}\} \\
&\wedge \exists z_i \exists z'_i \left[\begin{array}{l} \neg z_i =_i z'_i \\ \wedge \forall u_p [(z_i 2_{ip} u_p \vee z'_i 2_{ip} u_p) \rightarrow (y_i 2_{ip} u_p \wedge \neg x_i 2_{ip} u_p)] \\ \wedge \forall v_i [\rightarrow \exists t_p (x_i 2_{ip} t_p \wedge v_i 2_{ip} t_p)] \end{array} \right]
\end{aligned}$$

First suppose that $\mathcal{F} \models \varphi([a, b], [c, d])$. Then $[c, d]$ must *contain* every point *contained* by $[a, b]$, and there must be some point *contained* in $[c, d]$ which is not *contained* in $[a, b]$. So we have $c \leq a$, $b \leq d$ and $[a, b] \neq [c, d]$. (Note that here we are using the fact, which follows

| <i>Proved</i> | $=_p$ | $=_i$ | $<$ | 0_{ip} | 1_{ip} | 2_{ip} | 3_{ip} | 4_{ip} | 34_{ii} | 14_{ii} | 03_{ii} | 24_{ii} | 04_{ii} | 44_{ii} |
|--|-------|-------|-----|----------|----------|----------|----------|----------|-----------|-----------|-----------|-----------|-----------|-----------|
| $\{2_{ip}\} \cup \mathcal{J}^+$ | • | | | | | | | | | | | | | |
| $\{=_p, <, 0_{ip}, 1_{ip}\}$ | | • | | | | | | | | | | | | |
| $\{=_p, =_i, 1_{ip}, 14_{ii}\}$ | | | • | • | | | | | | | | | | |
| $\{=_p, 2_{ip}\} \cup \mathcal{J}^+$ | | | • | • | • | | • | • | | | | | | |
| $\{=_p, =_i, <, 3_{ip}, 4_{ip}, 03_{ii}\}$ | | | | • | • | • | | | • | • | | • | • | • |
| $\{=_p, <\} \cup \mathcal{J}^+$ | | | | • | • | • | • | • | | | | | | |
| $\{=_p, =_i, <, 0_{ip}, 1_{ip}\}$ | | | | | | | | | | • | | | | |
| $\{=_p, =_i, 2_{ip}, 04_{ii}\}$ | | | | | | | | | • | • | • | • | | • |
| $\{=_p, =_i, <, 04_{ii}\}$ | | | | | | | | | • | • | • | • | | • |
| <i>Symmetric</i> | $=_p$ | $=_i$ | $<$ | 0_{ip} | 1_{ip} | 2_{ip} | 3_{ip} | 4_{ip} | 34_{ii} | 14_{ii} | 03_{ii} | 24_{ii} | 04_{ii} | 44_{ii} |
| $\{=_p, <, 3_{ip}, 4_{ip}\}$ | | • | | | | | | | | | | | | |
| $\{=_p, =_i, 3_{ip}, 03_{ii}\}$ | | | • | | | | | • | | | | | | |
| $\{=_p, =_i, <, 0_{ip}, 1_{ip}, 14_{ii}\}$ | | | | | | • | • | • | • | | • | • | • | • |
| $\{=_p, =_i, <, 3_{ip}, 4_{ip}\}$ | | | | | | | | | | | • | | | |

Table 11: MIS(r), for each $r \in \mathfrak{R}^+$; upper part: sets for which we give an explicit construction; lower part: symmetric ones. - Class: Den.

from the density of the order, that $[a, b]$ contain at least one point. If this were not the case, the first conjunct of ϕ would be vacuously satisfied.) We need to show that $a \neq c$ and $d \neq b$. Suppose, to the contrary, that $a = c$. Then we must have that $b < d$. Let ψ denote the third main conjunct of φ and let $z_i = [e, f]$ and $z'_i = [g, h]$ be two intervals satisfying ψ . Then $[e, f] \neq [g, h]$ by the first conjunct of ψ , while $b \leq e < f \leq d$ and $b \leq g < h \leq d$. However now, taking $v_i = [b, d]$ falsifies the third conjunct of ψ , as $[b, d]$ contains all points contained in either $[e, f]$ or $[g, h]$ while not containing any point contained in $x_i = [a, b]$. This is a contradiction. Conversely, if $[a, b]04_{ii}[c, d]$, the first conjunct of φ holds trivially, while any point e such that $c < e < a$ or $b < e < d$ (which exists because the structure is dense) witnesses the z_p of the second conjunct. Taking $z_i = [c, a]$ and $z'_i = [b, d]$ witnesses the third conjunct. All that remains to be proved is the 44_{ii} -completeness of the sets $\{0_{ip}, 04_{ii}\}$, $\{1_{ip}, 04_{ii}\}$, $\{2_{ip}, <\}$, $\{3_{ip}, 04_{ii}\}$ and $\{4_{ip}, 04_{ii}\}$. But all of these sets have already been shown to define 34_{ii} over Den, and we know that $\{34_{ii}\}$ defines 44_{ii} over Lin. \square

4. INCOMPLETENESS RESULTS IN THE CLASS Den

We can now turn our attention to the maximal incomplete sets for relations in \mathfrak{R}^+ . Notice that for some $r \in \mathfrak{R}^+$ some r -incomplete set in the class Lin is also maximally r -incomplete in the class Den; nevertheless, all proofs in Part I must be revisited here, as they were designed to work not only in Lin but also in Dis, and therefore are not valid in Den.

Lemma 18. Tab. 11 is correct.

Proof. Let S be $\{2_{ip}\} \cup \mathcal{J}^+$: proving that it is $=_p$ -incomplete is almost immediate. Indeed, it suffices to take \mathbb{D} and \mathbb{D}' both equal to the subset of \mathbb{Q} of all points between 0 and 1, $\zeta = (\zeta_p, \zeta_i)$, where $\zeta_i = Id_i$ (the identical relation on intervals), $\zeta_p = \{(0, 1')\}$ plus the identical relation on points to have a surjective truth-preserving relation that breaks $=_p$.

Proving that $\{\underline{=}_p, <, 0_{ip}, 1_{ip}\}$ is $=_i$ -incomplete is equally easy: it suffices to take $\mathbb{D} = \mathbb{D}' = \mathbb{Q}$, $\zeta_p = Id_p$, and $\zeta_i = Id_i$ plus $\zeta_i([0, 2], [0', 1'])$. Assume, now, S to be $\{\underline{=}_p, =_i, 1_{ip}, 14_{ii}\}$; we need to prove its $<, 0_{ip}$ -incompleteness in the dense case. Let $\mathbb{D} = \mathbb{D}' = \mathbb{Q}$, and define $\zeta = \zeta_p \cup \zeta_i$ as follows: $(a, -a') \in \zeta_p$ for every $a \in \mathbb{Q}$, and $([a, b], [-a', -a' + |b' - a'|]) \in \zeta_i$ for every $[a, b] \in \mathbb{I}(\mathbb{Q})$, so that the length of every interval is preserved while their beginning point are reflected over 0; as this relation breaks both $<$ and 0_{ip} , the latter cannot be expressed in this language. Let now S be $\{\underline{=}_p, 2_{ip}\} \cup \mathcal{J}^+$: we can prove that it is $<, 0_{ip}, 1_{ip}, 3_{ip}, 4_{ip}$ -incomplete. To this end, it suffices to take, once again, \mathbb{D} and \mathbb{D}' both equal to the subset of \mathbb{Q} of all points between 0 and 1, $\zeta = (\zeta_p, \zeta_i)$, where $\zeta_i = Id_i$ (the identical relation on intervals), $\zeta_p = \{(0, 1'), (1, 0')\}$ plus the identical relation on every other point to have a surjective truth-preserving relation that breaks the relations under analysis. As for $S = \{\underline{=}_p, =_i, <, 3_{ip}, 4_{ip}, 03_{ii}\}$, we can prove its $0_{ip}, 1_{ip}, 2_{ip}, i$ -incompleteness, where $i \in \mathcal{J}^+ \setminus \{=_i, 03_{ii}\}$ by defining two \mathbb{Q} -based structures and a relation between them defined as the identity between points and as $\zeta_i([a, b]) = [a' - |b' - a'|, b']$, obtaining (as we did for the same set of relations on Lin , using, in that case, a pseudo-discrete structure) a relation that maps every interval to the interval with the same *ending* point but twice the length. In this way, all relations in S are respected. The m -incompleteness of $\{\underline{=}_p, <\} \cup \mathcal{J}^+$ for each $m \in \mathfrak{M}^+$ can be proved by taking again $\mathbb{D} = \mathbb{D}' = \mathbb{Q}$, $\zeta = (\zeta_p, \zeta_i)$, where $\zeta_i = Id_i$ and $\zeta_p(a) = a' + 1$ for every $a \in \mathbb{Q}$, which clearly respects all interval-interval relations, and both equality and relative ordering between points, but breaks every relation between points and intervals. When S is $\{\underline{=}_p, =_i, <, 0_{ip}, 1_{ip}\}$, we have to prove that it is 14_{ii} -incomplete. Consider two structures based on \mathbb{Q} , and let $\zeta = (\zeta_p, \zeta_i)$ be defined as $\zeta_p = Id_p$, and $\zeta_i = Id_i$ except for the interval $[-1, 0]$, which is mapped to $[-1, 1]$. When S is $\{\underline{=}_p, =_i, 2_{ip}, 04_{ii}\}$, we have to prove that it is r -incomplete in the dense case for $r \in \mathcal{J}^+ \setminus \{04_{ii}, =_i\}$. Consider two structures based on \mathbb{Q} , and let $\zeta = (\zeta_p, \zeta_i)$ be defined as $\zeta_p(a) = -a'$ for every point and $\zeta_i([a, b]) = [-b', -a']$ for every interval. Clearly, *containment* is respected for both sorts; nevertheless, all other interval-interval relations are broken. Finally, when S is $\{\underline{=}_p, =_i, <, 04_{ii}\}$, we have to prove that it is r -incomplete for $r \in \mathcal{J}^+ \setminus \{=_i, 04_{ii}\}$. Consider two structures based on \mathbb{Q} , and let $\zeta = (\zeta_p, \zeta_i)$ be defined as $\zeta_p = Id_p$, and $\zeta_i([a, b]) = [-b, -a]$; again, we respect *containment* between intervals, and the relative ordering between points is respected as well (since points are not affected by the construction), and we break every other interval-interval relation. \square

5. COMPLETENESS RESULTS IN THE CLASS Unb

The ability of fragments of our language to define relations when the underlying linear order is unbounded (but not necessarily discrete or dense) differs from the dense/discrete cases only slightly. Following the same schema, we now focus on the definability part, again, pointing out the differences with the linear case.

5.1. Definability in \mathcal{J}^+ and in \mathfrak{M}^+ . We begin our study of definability over the Class Unb by considering minimal definability of \mathcal{J}^+ relations by sets of \mathcal{J}^+ relations, and then of \mathfrak{M}^+ relations by sets of \mathfrak{M}^+ relations.

Lemma 19. Tab. 12 is correct.

| 14_{ii} | 34_{ii} | 24_{ii} | 04_{ii} | 44_{ii} |
|------------------------|------------------------|------------------------|------------------------|------------------------|
| $\{24_{ii}, 03_{ii}\}$ | $\{14_{ii}, 24_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ | $\{14_{ii}, 24_{ii}\}$ | $\{14_{ii}, 24_{ii}\}$ |
| $\{03_{ii}, 04_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ | $\{14_{ii}, 04_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ |
| $\{34_{ii}\}$ | $\{14_{ii}, 04_{ii}\}$ | $\{03_{ii}, 04_{ii}\}$ | $\{24_{ii}, 03_{ii}\}$ | $\{14_{ii}, 04_{ii}\}$ |
| $\{44_{ii}\}$ | $\{24_{ii}, 03_{ii}\}$ | $\{34_{ii}\}$ | $\{34_{ii}\}$ | $\{24_{ii}, 03_{ii}\}$ |
| | $\{24_{ii}, 04_{ii}\}$ | $\{44_{ii}\}$ | $\{44_{ii}\}$ | $\{24_{ii}, 04_{ii}\}$ |
| | $\{03_{ii}, 04_{ii}\}$ | | | $\{03_{ii}, 04_{ii}\}$ |
| | $\{44_{ii}\}$ | | | $\{34_{ii}\}$ |

Table 12: The spectrum of the $\text{mcs}(r)$, for each $r \in \mathcal{I}^+ \setminus \{=_{ii}\}$. - Class: Unb.

Proof. To prove that $\{44_{ii}\}$ is 14_{ii} -complete we use the same definition as was used in the dense case in Lemma 11:

$$x_i 14_{ii} y_i \leftrightarrow \forall z_i (z_i 44_{ii} x_i \leftrightarrow z_i 44_{ii} y_i) \wedge \exists z_i (x_i 44_{ii} z_i \wedge \neg(y_i 44_{ii} z_i)). \quad \{44_{ii}\}$$

The argument for the correctness of this definition is the same as in the dense case except that, when arguing that $[a, b] 14_{ii} [c, d]$ implies $\phi([a, b], [c, d])$, we appeal to unboundedness rather than density to justify the existence of the interval required by the second conjunct of the definition. As before, the 34_{ii} -completeness of $\{44_{ii}\}$ now follows via 34_{ii} -completeness of $\{14_{ii}, 44_{ii}\}$ over all linear orders. Now, since 34_{ii} defines all interval-interval relations over all linear orders, it follows that 44_{ii} is also complete with respect all other interval-interval relations over unbounded orders. To complete the proof, it suffices to show the 34_{ii} -completeness of $\{24_{ii}, 04_{ii}\}$. We do so by proving, as we did in Lemma 3.1 of Part I of this paper, that this set is able to express the weaker relation $34_{ii} \cup 44_{ii}$:

$$\begin{aligned} x_i 34_{ii} \cup 44_{ii} y_i \leftrightarrow & \exists z_i, k_i (x_i 04_{ii} z_i \wedge y_i 04_{ii} k_i \wedge z_i 24_{ii} k_i) \wedge & \{04_{ii}, 24_{ii}\} \\ & \exists z_i (x_i 04_{ii} z_i \wedge \neg(y_i 04_{ii} z_i)) \wedge \exists k_i (y_i 04_{ii} k_i \wedge \neg(x_i 04_{ii} k_i)) \wedge \\ & \neg(x_i 24_{ii} y_i) \wedge \neg(y_i 24_{ii} x_i) \wedge \neg(x_i 04_{ii} y_i) \wedge \neg(y_i 04_{ii} x_i). \end{aligned}$$

If $\mathcal{F} \models \varphi([a, b], [c, d])$ over an unbounded structure, we can eliminate all possibilities for the relationship between $[a, b]$ and $[c, d]$ other than 34_{ii} and 44_{ii} . Indeed, first observe that x_i and y_i cannot *overlap* nor *contain* each other. Next, if y_i *ends* before x_i or at its *beginning* point, it would be impossible to place z_i and k_i . Finally, if x_i *starts* or *finishes* y_i , or the other way around, we have a contradiction with the second or the third requirement of φ . This implies that x_i *meets* or is *before* y_i . Conversely, if $[a, b] 34_{ii} \vee 44_{ii} [c, d]$, then we can take $z_i = [e, d]$, where $e < a$, and $k_i = [a, f]$, where $d < f$, and the existence of e and f is guaranteed by the assumption of unboundedness. \square

5.2. Definability for the Relations $=_p, =_i, <$ in Unb.

Lemma 20. Tab. 13 is correct.

| $=_p$ | $=_i$ | $<$ |
|--------------|-----------------------|-----------------------|
| $\{<\}$ | $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}\}$ |
| $\{0_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{1_{ip}, 2_{ip}\}$ |
| $\{2_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ |
| $\{1_{ip}\}$ | $\{0_{ip}, 24_{ii}\}$ | $\{1_{ip}, 24_{ii}\}$ |
| $\{3_{ip}\}$ | $\{1_{ip}, 24_{ii}\}$ | $\{1_{ip}, 03_{ii}\}$ |
| $\{4_{ip}\}$ | $\{1_{ip}, 2_{ip}\}$ | $\{1_{ip}, 34_{ii}\}$ |
| | $\{1_{ip}, 3_{ip}\}$ | $\{1_{ip}, 04_{ii}\}$ |
| | $\{1_{ip}, 4_{ip}\}$ | $\{1_{ip}, 44_{ii}\}$ |
| | $\{2_{ip}, 3_{ip}\}$ | $\{2_{ip}, 3_{ip}\}$ |
| | $\{2_{ip}, 4_{ip}\}$ | $\{2_{ip}, 14_{ii}\}$ |
| | $\{3_{ip}, 24_{ii}\}$ | $\{2_{ip}, 24_{ii}\}$ |
| | $\{4_{ip}, 24_{ii}\}$ | $\{2_{ip}, 03_{ii}\}$ |
| | $\{14_{ii}\}$ | $\{2_{ip}, 34_{ii}\}$ |
| | $\{03_{ii}\}$ | $\{2_{ip}, 44_{ii}\}$ |
| | $\{34_{ii}\}$ | $\{3_{ip}, 14_{ii}\}$ |
| | $\{04_{ii}\}$ | $\{3_{ip}, 24_{ii}\}$ |
| | $\{44_{ii}\}$ | $\{3_{ip}, 34_{ii}\}$ |
| | | $\{3_{ip}, 04_{ii}\}$ |
| | | $\{3_{ip}, 44_{ii}\}$ |
| | | $\{4_{ip}\}$ |

Table 13: The spectrum of the $mcs(r)$, for each $r \in \{=_p, =_i, <\}$. - Class: Unb.

Proof. Starting with $=_p$, we have now that every mixed relation is $=_p$ -complete, as follows:

$$x_p =_p y_p \leftrightarrow \forall x_i (x_i m x_p \leftrightarrow x_i m y_p). \quad \{m\}, m \in \mathfrak{M}^+$$

Observe that this definition is the same that we have used in the dense case for 0_{ip} and 4_{ip} ; the difference is that now, because of the unboundedness hypothesis, the argument also works for $\{2_{ip}\}$, as we now prove. The left to right direction is immediate. We argue contrapositively for the other implication. Suppose that $x_p = a$ and $y_p = b$ are not equal. Without loss of generality we can assume that $x_p < y_p$. Since the underlying domain is unbounded, there must be a point $c < a$, and therefore, the interval $[c, b]$ is such that $[c, b] \not\subseteq_2 a$ but it is not the case that $[c, b] \subseteq_2 b$, falsifying the right-hand side. Notice that this argument does not work on dense structures that are left/right bounded, such as $[0, 1] \subset \mathbb{Q}$. Now, consider the relation $=_i$, and the following definitions:

$$x_i =_i y_i \leftrightarrow \begin{cases} \forall z_p (x_i 0_{ip} z_p \leftrightarrow y_i 0_{ip} z_p) \wedge \forall u_i (x_i 24_{ii} u_i \leftrightarrow y_i 24_{ii} u_i) & \{0_{ip}, 24_{ii}\} \\ \forall z_p (x_i 1_{ip} z_p \leftrightarrow y_i 1_{ip} z_p) \wedge \forall u_i (x_i 24_{ii} u_i \leftrightarrow y_i 24_{ii} u_i) & \{1_{ip}, 24_{ii}\} \\ \forall z_i (x_i 04_{ii} z_i \leftrightarrow y_i 04_{ii} z_i) \wedge \forall z_i (z_i 04_{ii} x_i \leftrightarrow z_i 04_{ii} y_i) & \{04_{ii}\} \\ \forall z_i (x_i 44_{ii} z_i \leftrightarrow y_i 44_{ii} z_i) \wedge \forall z_i (z_i 44_{ii} x_i \leftrightarrow z_i 44_{ii} y_i). & \{44_{ii}\} \end{cases}$$

First take the set $\{0_{ip}, 24_{ii}\}$. If $x_i =_i y_i$ we immediately have $\varphi(x_i, y_i)$. Conversely, suppose $x_i = [a, b] \neq_i [c, d] = y_i$, so either $a \neq c$ or $b \neq d$. In the first case the first conjunct of φ is falsified. Assume therefore that $a = c$ but that $b \neq d$ and, w.l.o.g., that $b < d$. Since the order is unbounded, there is a point $e > d$. Then $[c, d]24_{ii}[b, e]$ but $\neg([a, b]24_{ii}[b, e])$, falsifying the second conjunct. The argument in the case of $\{1_{ip}, 24_{ii}\}$ is virtually identical. The $=_i$ -completeness of $\{4_{ip}, 24_{ii}\}$ and $\{3_{ip}, 24_{ii}\}$ follows by symmetry (Lemma 10). Consider the set $\{04_{ii}\}$. If $x_i =_i y_i$ we immediately have $\varphi(x_i, y_i)$. Conversely, suppose that $[a, b] \neq_i [c, d]$, so either $a \neq c$ or $b \neq d$. Consider, w.l.o.g., the case $a \neq c$, specifically $a < c$. We then choose a point $e > \max\{b, d\}$, which does exist because the underlying domain is unbounded. Now, we have that $[c, d]04_{ii}[a, e]$, but it is not the case that $[a, b]04_{ii}[a, e]$, falsifying the right-hand side. The other case is treated with a similar argument. Let us now focus on $<$. Five new definitions are needed:

$$x_p < y_p \leftrightarrow \begin{cases} \forall x_i(x_i 0_{ip} y_p \rightarrow x_i 0_{ip} x_p) \wedge \exists x_i(x_i 0_{ip} x_p \wedge \neg(x_i 0_{ip} y_p)) & \{0_{ip}\} \\ \exists x_i y_i(x_i 1_{ip} x_p \wedge y_i 1_{ip} y_p \wedge x_i 24_{ii} y_i) & \{1_{ip}, 24_{ii}\} \\ \exists x_i y_i(x_i 2_{ip} x_p \wedge y_i 2_{ip} y_p \wedge \neg(x_i 2_{ip} y_p) \wedge \neg(y_i 2_{ip} x_p) \wedge x_i 24_{ii} y_i) & \{2_{ip}, 24_{ii}\} \\ \exists x_i y_i(x_i 2_{ip} x_p \wedge y_i 2_{ip} y_p \wedge \neg(x_i 2_{ip} y_p) \wedge x_i 14_{ii} y_i) & \{2_{ip}, 14_{ii}\} \\ \exists x_i y_i(x_i 2_{ip} x_p \wedge y_i 2_{ip} y_p \wedge \neg(x_i 2_{ip} y_p) \wedge \neg(y_i 2_{ip} x_p) \wedge \\ \exists z_i, t_i(z_i 44_{ii} y_i \wedge \neg(z_i 44_{ii} x_i) \wedge x_i 44_{ii} t_i \wedge \neg(y_i 44_{ii} t_i)) & \{2_{ip}, 44_{ii}\} \end{cases}$$

Consider, first the set $\{0_{ip}\}$. Suppose that $\mathcal{F} \models \varphi(a, b)$, and, for the sake of contradiction, that $a \geq b$. If $a = b$, then every interval $[c, d]$ such that $[c, d]0_{ip}a$ must be such that $[c, d]0_{ip}b$ as well, contradicting the second conjunct, and, if $b < a$ then, by unboundedness, there exists an interval $[a, c]$ such that $[a, c]0_{ip}b$ but it is not the case that $[a, c]0_{ip}a$, contradicting the first conjunct. On the other hand, suppose that $x_p = a < b$. Then if $[c, d]0_{ip}b$, we have $b < c$ and hence $a < c$, therefore $[c, d]0_{ip}a$, satisfying the first conjunct. For the sake of the second conjunct, consider any interval $[b, c]$: such an interval exists by unboundedness and clearly $[b, c]0_{ip}a$ while it is not the case that $[b, c]0_{ip}b$. Consider, now, the set $\{1_{ip}, 24_{ii}\}$. If $x_p = a < b = y_p$ then, by unboundedness, we may choose points c and d such that $a < b < c < d$. Then the intervals $x_i = [a, c]$ and $y_i = [b, d]$ witness φ . Conversely, if $\mathcal{F} \models \varphi(a, b)$, then there must be intervals $[a, c]$ and $[b, d]$ such that $[a, c]24_{ii}[b, d]$, i.e. such that $a < b < c < d$. So, in particular, $a < b$. As for $\{2_{ip}, 24_{ii}\}$, suppose that $\mathcal{F} \models \varphi(a, b)$, and, by contradiction, that $a \geq b$. If $a = b$, then for any interval *containing* a also *contains* b , and vice versa, making the first three conjuncts not simultaneously satisfiable. If, on the other hand, $b < a$, any interval *containing* b must *start* before a , and therefore no interval *containing* a can *overlap* an interval *containing* b , making $\varphi(a, b)$ false. Conversely, assume that $a < b$. By unboundedness we can take two intervals $[c, b]$ such that $c < a$ and $[a, d]$ such that $b < d$ to witness $\varphi(a, b)$, as we wanted. As for $\{2_{ip}, 14_{ii}\}$, suppose that $\mathcal{F} \models \varphi(a, b)$, and, by contradiction, that $a \geq b$. If $a = b$, then a and b are contained by the same intervals, making the first three conjuncts not simultaneously satisfiable. If, on the other hand, $b < a$, then any interval *containing* a but not b must *start* at b or after, and therefore it cannot *start* an interval containing b , falsifying $\varphi(a, b)$. Conversely, suppose that $a < b$: then, the intervals $[a, b]$ and $[a, c]$ for some $c > b$ witness $\varphi(a, b)$, as we wanted.

Finally, as for $\{2_{ip}, 44_{ii}\}$, suppose that $\mathcal{F} \models \varphi(a, b)$, and, by contradiction, that $a \geq b$. If $a = b$, then a and b are contained by the same intervals, making the first four conjuncts not simultaneously satisfiable. If, on the other hand, $b < a$, then any interval *containing* a but not b must *start* at b or after it, and every interval *containing* b but not a must *end* at a or before it; thus, z_i and t_i cannot be witnessed by any concrete interval, as z_i should *end* before y_i *starts* but not before x_i does, and, symmetrically, x_i should *end* before t_i *starts* but not before y_i does. Conversely, suppose that $a < b$: then, any interval of the type $[c, b]$ ($c < a$) serves as witness of x_i , and any interval of the type $[a, d]$ ($d > b$) serves as witness of y_i , while z_i and t_i are witnessed by any intervals $[e, c]$ ($e < c$) and $[d, f]$ ($d < f$), respectively. Such points c, d, e and f exist thanks to the unboundedness of the order. Every other definition is now straightforward. The cases for $\{2_{ip}, 34_{ii}\}$ follows since $\{34_{ii}\}$ defines 14_{ii} in the general case and the just-proved $<$ -completeness of $\{2_{ip}, 14_{ii}\}$. The cases for $\{1_{ip}, 04_{ii}\}$ and $\{1_{ip}, 44_{ii}\}$ follow since both $\{04_{ii}\}$ and $\{44_{ii}\}$ define $=_i$ in the unbounded case and both $\{1_{ip}, 04_{ii}, =_i\}$ and $\{1_{ip}, 44_{ii}, =_i\}$ define $<$ in the general case. The remaining cases, viz. $\{2_{ip}, 03_{ii}\}$, $\{3_{ip}, 24_{ii}\}$, $\{3_{ip}, 04_{ii}\}$, $\{3_{ip}, 44_{ii}\}$ and $\{4_{ip}\}$, now follow by symmetry and Lemma 10 from the already-proved cases for $\{2_{ip}, 14_{ii}\}$, $\{1_{ip}, 24_{ii}\}$, $\{1_{ip}, 04_{ii}\}$, $\{1_{ip}, 44_{ii}\}$ and $\{0_{ip}\}$, respectively. \square

5.3. Definability for \mathfrak{M}^+ -Relations in Unb. Because the relationship between definability of \mathfrak{M}^+ and \mathfrak{J}^+ over unbounded orders is somewhat involved, the corresponding results cannot be kept as neatly apart as e.g. in the dense case. We therefore first present and prove two lemmas (Lemmas 21 and 22) which collect some key definability results for \mathfrak{M}^+ and \mathfrak{J}^+ , respectively. These results are then used to prove the full collection of minimal definability results for \mathfrak{M}^+ and \mathfrak{J}^+ (Lemmas 23 and 24).

Lemma 21. Over unbounded linear orders the sets $\{2_{ip}, 04_{ii}, <\}$ and $\{1_{ip}, 24_{ii}\}$ are 0_{ip} -complete, the set $\{0_{ip}\}$ is 1_{ip} -complete, while the sets $\{1_{ip}, 24_{ii}\}$, $\{1_{ip}, 04_{ii}\}$, $\{0_{ip}, 24_{ii}\}$ and $\{0_{ip}, 04_{ii}\}$ are 2_{ip} -complete.

Proof. Consider the following definition of 0_{ip} in terms of $2_{ip}, 04_{ii}$ and $<$:

$$x_i 0_{ip} y_p \leftrightarrow \exists z_i (x_i 04_{ii} z_i \wedge \forall k_p (z_i 2_{ip} k_p \rightarrow y_p < k_p)). \quad \{2_{ip}, 04_{ii}, <\}$$

Proceeding as always, suppose that $\mathcal{F} \models \varphi([a, b], c)$. Let $z_i = [d, e]$ the interval witnessing φ , so that $d < a < b < e$. Since $z_i 2_{ip} a$ by definition it follows that $c < a$, i.e. that $[a, b] 0_{ip} c$. Conversely, suppose that $[a, b] 0_{ip} c$. Since the underlying order is unbounded there is a point d such that $b < d$. Then the interval $[c, d]$ witnesses the definition. As for 1_{ip} , we show that $\{0_{ip}\}$ is 1_{ip} -complete by noticing that $\{0_{ip}\}$ is $<, =_p$ -complete and by using the following, straightforward definition:

$$x_i 1_{ip} y_p \leftrightarrow \forall z_p (x_i 0_{ip} z_p \rightarrow z_p < y_p) \wedge \forall z_p (\neg(x_i 0_{ip} z_p) \rightarrow (y_p < z_p \vee y_p =_p z_p)). \quad \{0_{ip}\}$$

Finally, consider the following two new definitions for 2_{ip} :

$$x_i 2_{ip} y_p \leftrightarrow \begin{cases} \exists z_i (z_i 1_{ip} y_p \wedge x_i 2_{ii} z_i) & \{1_{ip}, 2_{ii}\} \\ \neg(x_i 1_{ip} y_p) \wedge \exists z_i (z_i 1_{ip} y_p \wedge \forall k_i (x_i 0_{ii} k_i \rightarrow z_i 0_{ii} k_i)). & \{1_{ip}, 0_{ii}\} \end{cases}$$

The first one, namely $\{1_{ip}, 2_{ii}\}$, is straightforward. As for $\{1_{ip}, 0_{ii}\}$, suppose first that $\mathcal{F} \models \varphi([a, b], c)$. If $c < a$, then we have the interval $[c, a]$ that contradicts the definition; c cannot be a , and if $c \geq b$, then every interval $[c, d]$, where $d > c \geq b$ exists by hypothesis, again, leads to a contradiction. Thus, $a < c < b$. If, on the other hand, $a < c < b$ then $z_i = [c, b]$ witnesses the existential quantifier in the definition. It remains to be proved the 0_{ip} -completeness of $\{1_{ip}, 2_{ii}\}$ and the 2_{ip} -completeness of $\{0_{ip}, 2_{ii}\}$ and $\{0_{ip}, 0_{ii}\}$. These do not require new definitions. Indeed, the 0_{ip} -completeness of $\{1_{ip}, 2_{ii}\}$ follows since $\{1_{ip}, 2_{ii}\}$ defines 2_{ip} over unbounded orders (as we proved above) and $\{1_{ip}, 2_{ip}\}$ defines 0_{ip} over all linear orders. That $\{0_{ip}, 2_{ii}\}$ and $\{0_{ip}, 0_{ii}\}$ are 2_{ip} -complete follows since $\{0_{ip}\}$ defines 1_{ip} over unbounded order while, as we have just proved, $\{1_{ip}, 2_{ii}\}$ and $\{1_{ip}, 0_{ii}\}$ both define 2_{ip} over these order. \square

Lemma 22. The sets $\{0_{ip}, 2_{ii}\}$, $\{1_{ip}, 2_{ii}\}$, $\{0_{ip}, 0_{ii}\}$, $\{1_{ip}, 0_{ii}\}$, $\{4_{ip}, 2_{ii}\}$, $\{3_{ip}, 2_{ii}\}$, $\{4_{ip}, 0_{ii}\}$, $\{3_{ip}, 0_{ii}\}$ are 3_{ii} -complete over unbounded linear orders.

Proof. No new definitions are required, as all these definability results follow from previous results. Indeed, we begin by recalling that $\{0_{ip}, 2_{ip}\}$ defines 3_{ii} over all linear orders. From this the cases of $\{0_{ip}, 2_{ii}\}$ and $\{0_{ip}, 0_{ii}\}$ follow via the fact that both of them define 2_{ip} over unbounded orders (Lemma 21). The cases for $\{1_{ip}, 2_{ii}\}$ and $\{1_{ip}, 0_{ii}\}$ follow similarly since both define 2_{ip} over unbounded orders (Lemma 21) and $\{1_{ip}, 2_{ip}\}$ defines 3_{ii} over all linear orders. Lastly, the cases for $\{4_{ip}, 2_{ii}\}$, $\{3_{ip}, 2_{ii}\}$, $\{4_{ip}, 0_{ii}\}$, $\{3_{ip}, 0_{ii}\}$ follow by symmetry (Lemma 10) from the previous ones. \square

Lemma 23. Tab. 14 is correct.

Proof. No new explicit definitions are needed. We focus first on 0_{ip} -completeness. The cases for $\{2_{ip}, 0_{ii}, <\}$ and $\{1_{ip}, 2_{ii}\}$ were already dealt with in Lemma 21. In the cases $\{1_{ip}, 0_{ii}\}$ and $\{1_{ip}, 4_{ii}\}$ we note that both $\{0_{ii}\}$ and $\{4_{ii}\}$ define $=_i$ in the unbounded case, and that both $\{1_{ip}, 0_{ii}, =_i\}$ and $\{1_{ip}, 4_{ii}, =_i\}$ are 0_{ip} -complete. For $\{2_{ip}, 1_{ii}\}$, note that $\{2_{ip}, 1_{ii}\}$ defines $<$ in the unbounded case and that $\{2_{ip}, 1_{ii}, <\}$ defines 0_{ip} . The case for $\{2_{ip}, 2_{ii}, 0_{ii}\}$ follows from the facts that, in the unbounded case, $\{2_{ip}, 2_{ii}\}$ defines $<$ and $\{2_{ip}, 0_{ii}, <\}$ defines 0_{ip} , as already proved. The case for $\{2_{ip}, 4_{ii}\}$ now follows from the latter via the fact that $\{4_{ii}\}$ is 2_{ii} -complete and $\{4_{ii}\}$ is 0_{ii} -complete in the unbounded case. For $\{2_{ip}, 0_{3ii}\}$, observe that $\{2_{ip}, 0_{3ii}\}$ defines $<$ in the unbounded case and that $\{2_{ip}, 0_{3ii}, <\}$ is 0_{ip} -complete. To see that $\{3_{ip}, 2_{ii}\}$ is 0_{ip} -complete, one may use the facts that $\{1_{ip}, 2_{ii}\}$ defines 3_{ii} in the unbounded case (Lemma 22), hence, by symmetry (Lemma 10), $\{3_{ip}, 2_{ii}\}$ defines 3_{ii} , and that $\{3_{ip}, 3_{ii}\}$ defines 0_{ip} in the general case. For $\{3_{ip}, 0_{ii}\}$, note that $\{3_{ip}, 0_{ii}\}$ defines 3_{ii} (Lemma 22) in the unbounded case, and then again that $\{3_{ip}, 3_{ii}\}$ defines 0_{ip} in the general case. The cases for $\{3_{ip}, 2_{ii}\}$, $\{3_{ip}, 0_{ii}\}$ and $\{3_{ip}, 4_{ii}\}$ all follow from the fact that $\{3_{ip}, 3_{ii}\}$ defines 0_{ip} , via the fact that each of these three sets defines 3_{ii} over Unb : in the first instance $\{1_{ip}, 2_{ii}\}$ defines 3_{ii} (Lemma 22), hence, by symmetry (Lemma 10) so does $\{3_{ip}, 2_{ii}\}$; in the second instance we have

| 0_{ip} | 1_{ip} | 2_{ip} |
|--------------------------------|--------------------------------|-----------------------|
| $\{1_{ip}, 2_{ip}\}$ | $\{0_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ |
| $\{1_{ip}, 3_{ip}\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ |
| $\{1_{ip}, 4_{ip}\}$ | $\{2_{ip}, 4_{ip}\}$ | $\{0_{ip}, 24_{ii}\}$ |
| $\{1_{ip}, 24_{ii}\}$ | $\{2_{ip}, 14_{ii}\}$ | $\{0_{ip}, 44_{ii}\}$ |
| $\{1_{ip}, 03_{ii}\}$ | $\{2_{ip}, 24_{ii}, 04_{ii}\}$ | $\{0_{ip}, 03_{ii}\}$ |
| $\{1_{ip}, 34_{ii}\}$ | $\{2_{ip}, 44_{ii}\}$ | $\{0_{ip}, 34_{ii}\}$ |
| $\{1_{ip}, 04_{ii}\}$ | $\{2_{ip}, 03_{ii}\}$ | $\{0_{ip}, 04_{ii}\}$ |
| $\{1_{ip}, 44_{ii}\}$ | $\{2_{ip}, 34_{ii}\}$ | $\{1_{ip}, 3_{ip}\}$ |
| $\{1_{ip}, <\}$ | $\{2_{ip}, 04_{ii}, <\}$ | $\{1_{ip}, 4_{ip}\}$ |
| $\{2_{ip}, 3_{ip}\}$ | $\{3_{ip}, 14_{ii}\}$ | $\{1_{ip}, 24_{ii}\}$ |
| $\{2_{ip}, 4_{ip}\}$ | $\{3_{ip}, 24_{ii}\}$ | $\{1_{ip}, 44_{ii}\}$ |
| $\{2_{ip}, 14_{ii}\}$ | $\{3_{ip}, 04_{ii}\}$ | $\{1_{ip}, 03_{ii}\}$ |
| $\{2_{ip}, 24_{ii}, 04_{ii}\}$ | $\{3_{ip}, 44_{ii}\}$ | $\{1_{ip}, 34_{ii}\}$ |
| $\{2_{ip}, 44_{ii}\}$ | $\{3_{ip}, 34_{ii}\}$ | $\{1_{ip}, 04_{ii}\}$ |
| $\{2_{ip}, 03_{ii}\}$ | $\{4_{ip}, 14_{ii}\}$ | $\{3_{ip}, 14_{ii}\}$ |
| $\{2_{ip}, 34_{ii}\}$ | $\{4_{ip}, 24_{ii}\}$ | $\{3_{ip}, 24_{ii}\}$ |
| $\{2_{ip}, 04_{ii}, <\}$ | $\{4_{ip}, 04_{ii}\}$ | $\{3_{ip}, 04_{ii}\}$ |
| $\{3_{ip}, 14_{ii}\}$ | $\{4_{ip}, 44_{ii}\}$ | $\{3_{ip}, 44_{ii}\}$ |
| $\{3_{ip}, 24_{ii}\}$ | $\{4_{ip}, 34_{ii}\}$ | $\{3_{ip}, 34_{ii}\}$ |
| $\{3_{ip}, 04_{ii}\}$ | | $\{4_{ip}, 14_{ii}\}$ |
| $\{3_{ip}, 44_{ii}\}$ | | $\{4_{ip}, 24_{ii}\}$ |
| $\{3_{ip}, 34_{ii}\}$ | | $\{4_{ip}, 04_{ii}\}$ |
| $\{4_{ip}, 14_{ii}\}$ | | $\{4_{ip}, 44_{ii}\}$ |
| $\{4_{ip}, 24_{ii}\}$ | | $\{4_{ip}, 34_{ii}\}$ |
| $\{4_{ip}, 04_{ii}\}$ | | |
| $\{4_{ip}, 44_{ii}\}$ | | |
| $\{4_{ip}, 34_{ii}\}$ | | |

Table 14: The spectrum of the $mcs(r)$, for each $r \in \mathfrak{M}^+$. - Class: Unb.

that $\{3_{ip}, 04_{ii}\}$ defines 34_{ii} (Lemma 22); in the third one we have that 44_{ii} defines, directly, 34_{ii} in the unbounded case. The set $\{4_{ip}, 14_{ii}\}$ is 0_{ip} -complete since $\{4_{ip}\}$ defines $<$ in the unbounded case and $\{4_{ip}, 14_{ii}, <\}$ defines 0_{ip} . The 0_{ip} -completeness of $\{4_{ip}, 34_{ii}\}$ follows from that of $\{4_{ip}, 14_{ii}\}$ since 34_{ii} defines all interval-interval relations over all linear orders. From this, in turn, follow the cases for $\{4_{ip}, 24_{ii}\}$ and $\{4_{ip}, 04_{ii}\}$ since both these sets define 34_{ii} over unbounded orders (Lemma 21). Lastly the case for $\{4_{ip}, 44_{ii}\}$ also follows from the 0_{ip} -completeness of $\{4_{ip}, 14_{ii}\}$ since 44_{ii} defines 14_{ii} over unbounded orders. We now turn our attention to 1_{ip} -completeness. The case of $\{0_{ip}\}$ was already treated in Lemma 21. All remaining cases follow immediately from this, as the sets involved have already been shown to define 0_{ip} over unbounded orders. Now focusing on 2_{ip} -completeness, the cases for $\{1_{ip}, 24_{ii}\}$, $\{1_{ip}, 04_{ii}\}$, $\{0_{ip}, 24_{ii}\}$ and $\{0_{ip}, 04_{ii}\}$ were already considered in Lemma 21. The cases for $\{0_{ip}, 44_{ii}\}$, $\{0_{ip}, 03_{ii}\}$ and $\{0_{ip}, 34_{ii}\}$ all follow from the 2_{ip} -completeness of $\{0_{ip}, 24_{ii}\}$ via, respectively, the facts that $\{44_{ii}\}$ defines 24_{ii} over unbounded orders, $\{0_{ip}, 03_{ii}\}$ defines 24_{ii} over linear orders, and $\{34_{ii}\}$ defines 24_{ii} over unbounded orders. The 2_{ip} -completeness of $\{1_{ip}, 44_{ii}\}$ follows from that of $\{1_{ip}, 24_{ii}\}$, since $\{44_{ii}\}$ defines 24_{ii} over

| 14_{ii} | 34_{ii} | 24_{ii} | 04_{ii} | 44_{ii} |
|--------------------------|--------------------------|------------------------|------------------------|--------------------------|
| $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}, 2_{ip}\}$ |
| $\{0_{ip}, 3_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ | $\{0_{ip}, 3_{ip}\}$ |
| $\{0_{ip}, 4_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ | $\{0_{ip}, 4_{ip}\}$ |
| $\{0_{ip}, 24_{ii}\}$ | $\{0_{ip}, 24_{ii}\}$ | $\{0_{ip}, 03_{ii}\}$ | $\{0_{ip}, 24_{ii}\}$ | $\{0_{ip}, 24_{ii}\}$ |
| $\{0_{ip}, 03_{ii}\}$ | $\{0_{ip}, 03_{ii}\}$ | $\{0_{ip}, 04_{ii}\}$ | $\{0_{ip}, 03_{ii}\}$ | $\{0_{ip}, 03_{ii}\}$ |
| $\{0_{ip}, 04_{ii}\}$ | $\{0_{ip}, 04_{ii}\}$ | $\{1_{ip}, 2_{ip}\}$ | $\{1_{ip}, 2_{ip}\}$ | $\{0_{ip}, 04_{ii}\}$ |
| $\{1_{ip}, 2_{ip}\}$ | $\{1_{ip}, 2_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ | $\{1_{ip}, 2_{ip}\}$ |
| $\{1_{ip}, 3_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ | $\{1_{ip}, 4_{ip}\}$ | $\{1_{ip}, 4_{ip}\}$ | $\{1_{ip}, 3_{ip}\}$ |
| $\{1_{ip}, 4_{ip}\}$ | $\{1_{ip}, 4_{ip}\}$ | $\{1_{ip}, 03_{ii}\}$ | $\{1_{ip}, 24_{ii}\}$ | $\{1_{ip}, 4_{ip}\}$ |
| $\{1_{ip}, 24_{ii}\}$ | $\{1_{ip}, 24_{ii}\}$ | $\{1_{ip}, 04_{ii}\}$ | $\{1_{ip}, 03_{ii}\}$ | $\{1_{ip}, 24_{ii}\}$ |
| $\{1_{ip}, 03_{ii}\}$ | $\{1_{ip}, 03_{ii}\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{1_{ip}, 03_{ii}\}$ |
| $\{1_{ip}, 04_{ii}\}$ | $\{1_{ip}, 04_{ii}\}$ | $\{2_{ip}, 4_{ip}\}$ | $\{2_{ip}, 4_{ip}\}$ | $\{1_{ip}, 04_{ii}\}$ |
| $\{2_{ip}, 3_{ip}\}$ | $\{2_{ip}, 3_{ip}\}$ | $\{2_{ip}, 14_{ii}\}$ | $\{2_{ip}, 14_{ii}\}$ | $\{2_{ip}, 3_{ip}\}$ |
| $\{2_{ip}, 4_{ip}\}$ | $\{2_{ip}, 4_{ip}\}$ | $\{2_{ip}, 03_{ii}\}$ | $\{2_{ip}, 03_{ii}\}$ | $\{2_{ip}, 4_{ip}\}$ |
| $\{24_{ii}, 04_{ii}\}$ | $\{2_{ip}, 14_{ii}\}$ | $\{2_{ip}, <\}$ | $\{3_{ip}, 14_{ii}\}$ | $\{2_{ip}, 14_{ii}\}$ |
| $\{2_{ip}, 03_{ii}\}$ | $\{2_{ip}, 03_{ii}\}$ | $\{3_{ip}, 14_{ii}\}$ | $\{4_{ip}, 14_{ii}\}$ | $\{2_{ip}, 03_{ii}\}$ |
| $\{2_{ip}, 04_{ii}, <\}$ | $\{2_{ip}, 04_{ii}, <\}$ | $\{4_{ip}, 14_{ii}\}$ | $\{14_{ii}, 24_{ii}\}$ | $\{2_{ip}, 04_{ii}, <\}$ |
| $\{24_{ii}, 03_{ii}\}$ | $\{3_{ip}, 14_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ | $\{3_{ip}, 14_{ii}\}$ |
| $\{03_{ii}, 04_{ii}\}$ | $\{4_{ip}, 14_{ii}\}$ | $\{14_{ii}, 04_{ii}\}$ | $\{24_{ii}, 03_{ii}\}$ | $\{4_{ip}, 14_{ii}\}$ |
| $\{34_{ii}\}$ | $\{14_{ii}, 24_{ii}\}$ | $\{03_{ii}, 04_{ii}\}$ | $\{34_{ii}\}$ | $\{14_{ii}, 24_{ii}\}$ |
| $\{44_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ | $\{34_{ii}\}$ | $\{44_{ii}\}$ | $\{14_{ii}, 03_{ii}\}$ |
| $\{3_{ip}, 24_{ii}\}$ | $\{14_{ii}, 04_{ii}\}$ | $\{44_{ii}\}$ | $\{3_{ip}, 24_{ii}\}$ | $\{14_{ii}, 04_{ii}\}$ |
| $\{4_{ip}, 24_{ii}\}$ | $\{24_{ii}, 03_{ii}\}$ | $\{3_{ip}, 04_{ii}\}$ | $\{4_{ip}, 24_{ii}\}$ | $\{24_{ii}, 03_{ii}\}$ |
| $\{3_{ip}, 04_{ii}\}$ | $\{24_{ii}, 04_{ii}\}$ | $\{4_{ip}, 04_{ii}\}$ | | $\{24_{ii}, 04_{ii}\}$ |
| $\{4_{ip}, 04_{ii}\}$ | $\{03_{ii}, 04_{ii}\}$ | | | $\{03_{ii}, 04_{ii}\}$ |
| | $\{44_{ii}\}$ | | | $\{34_{ii}\}$ |
| | $\{1_{ip}, 24_{ii}\}$ | | | $\{3_{ip}, 24_{ii}\}$ |
| | $\{3_{ip}, 24_{ii}\}$ | | | $\{4_{ip}, 24_{ii}\}$ |
| | $\{4_{ip}, 24_{ii}\}$ | | | $\{3_{ip}, 04_{ii}\}$ |
| | $\{3_{ip}, 04_{ii}\}$ | | | $\{4_{ip}, 04_{ii}\}$ |
| | $\{4_{ip}, 04_{ii}\}$ | | | |

Table 15: The spectrum of the $\text{mcs}(r)$, for each $r \in \mathcal{J}^+ \setminus \{=i\}$. - Class: Unb.

unbounded orders. The cases for $\{3_{ip}, 24_{ii}\}$ and $\{3_{ip}, 44_{ii}\}$ follow from the fact that both define 0_{ip} over unbounded orders and the 2_{ip} completeness of, respectively, $\{0_{ip}, 24_{ii}\}$ and $\{0_{ip}, 44_{ii}\}$. The 2_{ip} completeness of $\{4_{ip}, 14_{ii}\}$ follows by symmetry (Lemma 10) from that of $\{0_{ip}, 03_{ii}\}$. To see that $\{4_{ip}, 24_{ii}\}$, $\{4_{ip}, 04_{ii}\}$, $\{4_{ip}, 44_{ii}\}$ and $\{4_{ip}, 34_{ii}\}$ we note that each of these sets defines 0_{ip} over unbounded orders and than appeal to the 2_{ip} -completeness of the sets $\{0_{ip}, 24_{ii}\}$, $\{0_{ip}, 04_{ii}\}$, $\{0_{ip}, 44_{ii}\}$ and $\{0_{ip}, 34_{ii}\}$, respectively. \square

5.4. Definability for Relations in $\mathcal{J}^+ \setminus \{=i\}$ in Unb.

Lemma 24. Tab. 15 is correct.

Proof. We begin by noticing that every new 14_{ii} -, 24_{ii} -, 04_{ii} - or 44_{ii} -complete set also appears as a 34_{ii} -complete set. Since 34_{ii} defines every other interval-interval relation over linear orders, it is therefore sufficient to justify the 34_{ii} -completeness of all new 34_{ii} -complete sets.

| <i>Proved</i> | $=_p$ | $=_i$ | $<$ | 0_{ip} | 1_{ip} | 2_{ip} | 3_{ip} | 4_{ip} | 34_{ii} | 14_{ii} | 03_{ii} | 24_{ii} | 04_{ii} | 44_{ii} |
|--|-------|-------|-----|----------|----------|----------|----------|----------|-----------|-----------|-----------|-----------|-----------|-----------|
| \mathfrak{J}^+ | • | | | | | | | | | | | | | |
| $\{=_p, <, 0_{ip}, 1_{ip}\}$ | | • | | | | | | | | | | | | |
| $\{=_p, <, 2_{ip}, 24_{ii}\}$ | | • | | | | | | | | | | | | |
| $\{=_p, =_i, 1_{ip}, 14_{ii}\}$ | | | • | • | | | | | | | | | | |
| $\{=_p\} \cup \mathfrak{J}^+$ | | | • | | | | | | | | | | | |
| $\{=_p, <\} \cup \mathfrak{J}^+$ | | | | • | • | • | • | • | | | | | | |
| $\{=_p, =_i, 2_{ip}, 04_{ii}\}$ | | | • | • | • | | • | • | • | • | • | • | | • |
| $\{=_p, =_i, <, 3_{ip}, 4_{ip}, 03_{ii}\}$ | | | | • | • | • | | | • | • | | • | • | • |
| $\{=_p, =_i, <, 2_{ip}, 24_{ii}\}$ | | | | • | • | | • | • | • | • | | | • | • |
| $\{=_p, =_i, <, 0_{ip}, 1_{ip}\}$ | | | | | | | | | | • | | | | |
| $\{=_p, =_i, <, 04_{ii}\}$ | | | | | | | | | • | • | • | • | | • |
| <i>Symmetric</i> | $=_p$ | $=_i$ | $<$ | 0_{ip} | 1_{ip} | 2_{ip} | 3_{ip} | 4_{ip} | 34_{ii} | 14_{ii} | 03_{ii} | 24_{ii} | 04_{ii} | 44_{ii} |
| $\{=_p, <, 3_{ip}, 4_{ip}\}$ | | • | | | | | | | | | | | | |
| $\{=_p, =_i, 3_{ip}, 03_{ii}\}$ | | | • | | | | | • | | | | | | |
| $\{=_p, =_i, <, 0_{ip}, 1_{ip}, 14_{ii}\}$ | | | | | | • | • | • | • | | • | • | • | • |
| $\{=_p, =_i, <, 3_{ip}, 4_{ip}\}$ | | | | | | | | | | | • | | | |

Table 16: MIS(r), for each $r \in \mathfrak{R}^+$; upper part: sets for which we give an explicit construction; lower part: symmetric ones. - Class: Unb.

This has already been done for $\{24_{ii}, 04_{ii}\}$ and $\{44_{ii}\}$ in Lemma 19, and for $\{0_{ip}, 24_{ii}\}$, $\{1_{ip}, 24_{ii}\}$, $\{0_{ip}, 04_{ii}\}$, $\{1_{ip}, 04_{ii}\}$, $\{4_{ip}, 24_{ii}\}$, $\{3_{ip}, 24_{ii}\}$, $\{4_{ip}, 04_{ii}\}$, $\{3_{ip}, 04_{ii}\}$ in Lemma 22. The 34_{ii} -completeness of $\{1_{ip}, 24_{ii}\}$ follows from the 2_{ip} -completeness of $\{1_{ip}, 24_{ii}\}$ over unbounded orders and the 34_{ii} -completeness of $\{1_{ip}, 2_{ip}\}$ over linear orders. The case for $\{2_{ip}, 04_{ii}, <\}$ follows since this set defines 1_{ip} over unbounded orders and $\{1_{ip}, 2_{ip}\}$ is 34_{ii} -complete over linear orders. Lastly, the 34_{ii} -completeness of $\{4_{ip}, 14_{ii}\}$ is a consequence of the 1_{ip} - and 2_{ip} -completeness of this set over unbounded order and the 34_{ii} -completeness of $\{1_{ip}, 2_{ip}\}$ over linear orders. \square

6. INCOMPLETENESS RESULTS IN THE CLASS Unb

We can now turn our attention to the maximal incomplete sets for relations in \mathfrak{R}^+ . Notice that for some $r \in \mathfrak{R}^+$, some r -incomplete set in the class Den is also maximally r -incomplete in the class Unb, and it has been proven so by means of a dense unbounded counterexample; in these cases, we can borrow the same argument unchanged.

Lemma 25. Tab. 16 is correct.

Proof. Let S be \mathfrak{J}^+ : proving that it is $=_p$ -incomplete is very easy. Indeed, it suffices to take $\mathbb{D} = \mathbb{D}' = \mathbb{Q}$, $\zeta = (\zeta_p, \zeta_i)$, where $\zeta_i = Id_i$ (the identical relation on intervals), $\zeta_p = \{(0, 1')\}$ plus the identical relation on points to have a surjective truth-preserving relation that breaks $=_p$. The $=_i$ -incompleteness of $\{=_p, <, 0_{ip}, 1_{ip}\}$ is justified with the same argument used in the dense case (which was based on \mathbb{Q}). We can then prove that also $\{=_p, <, 2_{ip}, 24_{ii}\}$ is $=_i$ -incomplete, by taking $\mathbb{D} = \mathbb{D}' = \mathbb{Z}$, $\zeta_p = Id_p$, and $\zeta_i = Id_i$ plus $\zeta_i([1, 2], [0', 1'])$. For the $<$, 0_{ip} -incompleteness of $\{=_i, =_p, 1_{ip}, 14_{ii}\}$ we can recycle the argument used for

the dense case (again, based on \mathbb{Q}). The $<$ -incompleteness of $\{\underline{=}_p\} \cup \mathcal{I}^+$ can be proved by taking $\mathbb{D} = \mathbb{D}' = \mathbb{Q}$, $\zeta = (\zeta_p, \zeta_i)$, where $\zeta_i = Id_i$ and $\zeta_p(a) = -a'$ for every $a \in \mathbb{Q}$, which clearly respects all interval-interval relations and the equality between points. As for proving that $\{\underline{=}_p, <\} \cup \mathcal{I}^+$ is m -incomplete for each $m \in \mathfrak{M}^+$ we can recycle the same argument as in the dense case, as it was based on the set \mathbb{Q} . When S is $\{\underline{=}_p, =_i, 2_{ip}, 04_{ii}\}$, we have to prove that it is $<, 0_{ip}, 1_{ip}, i$ -incomplete, where $i \in \mathcal{I}^+ \setminus \{04_{ii}, =_i\}$. Consider two structures based on \mathbb{Q} , and let $\zeta = (\zeta_p, \zeta_i)$ be defined as $\zeta_p(a) = -a'$ for every point and $\zeta_i([a, b]) = [-b', -a']$ for every interval. Clearly, *containment* is respected for both sorts; nevertheless, $<, 0_{ip}$ and 1_{ip} and all interval-interval relations, except 04_{ii} , are broken. Once again, we have already proved that $\{\underline{=}_i, =_p, <, 3_{ip}, 4_{ip}, 03_{ii}\}$ is $0_{ip}, 1_{ip}, 2_{ip}, i$ -incomplete, where $i \in \mathcal{I}^+ \setminus \{03_{ii}, =_i\}$ when we were treating the dense case, and the same holds for the m, i -incompleteness of $\{\underline{=}_p, <\} \cup \mathcal{I}^+$, where $i \in \mathcal{I}^+$ and $m \in \mathfrak{M}^+$. Consider now the $0_{ip}, 1_{ip}, i$ -incompleteness of $\{\underline{=}_p, =_i, <, 2_{ip}, 24_{ii}\}$, where $i \in \mathcal{I}^+ \setminus \{24_{ii}, =_i\}$. Take $\mathbb{D} = \mathbb{D}' = \mathbb{Z}$, $\zeta_p = Id_p$ and $\zeta_i([a, a+1]) = [a'+1, a'+2]$ plus the identical relation over every other interval; since the only intervals affected by ζ are unitary, the relation 24_{ii} cannot be broken, and since such interval do not have internal points, the relation 2_{ip} cannot be broken either. Once more, the 14_{ii} -incompleteness of $\{\underline{=}_i, =_p, <, 0_{ip}, 1_{ip}\}$ comes directly from the dense case, and the same holds for the r -incompleteness of $\{\underline{=}_i, =_p, <, 04_{ii}\}$ for $r \in \mathcal{I}^+ \setminus \{04_{ii}\}$, which concludes the proof. \square

7. HARVEST: THE COMPLETE PICTURE FOR Den AND Unb

We are now capable to identify all expressively different subsets of \mathfrak{R}^+ under the hypotheses of linearity+density and linearity+unboundedness. Unlike Part I, we limit ourselves to list the maximally incomplete sets and the minimally complete sets for each of the two cases in the full language only.

Theorem 26. If a set of relations is listed:

- as $\text{mcs}(\mathfrak{R}^+)$ in Tab. 17, left column (resp., right column), then it is minimally \mathfrak{R}^+ -complete (resp., maximally \mathfrak{R}^+ -incomplete) in the class of all dense linearly ordered sets.
- as $\text{mcs}(\mathfrak{R}^+)$ in Tab. 18, left column (resp., right column), then it is minimally \mathfrak{R}^+ -complete (resp., maximally \mathfrak{R}^+ -incomplete) in the class of all unbounded linearly ordered sets.

8. CONCLUSIONS

We considered here the two-sorted first-order temporal languages that includes relations between intervals, points, and inter-sort, and we treated equality between points and between intervals as any other relation, with no special role. Under four different assumptions on the underlying structure, namely, linearity only, linearity+discreteness, linearity+density, and linearity+unboundedness, we asked the question: which relation can be first-order defined by which subset of all relations? As a result, we identified all possible inter-definability between relations, all minimally complete, and all maximally incomplete subsets of relations. These inter-definability results allow one to effectively compute all expressively different

| \mathfrak{R}^+ | |
|-----------------------|--|
| mcs | MIS |
| $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}, 1_{ip}, 14_{ii}, =_i, =_p, <\}$ |
| $\{0_{ip}, 3_{ip}\}$ | $\{2_{ip}, 14_{ii}, 24_{ii}, 03_{ii}, 34_{ii}, 04_{ii}, 44_{ii}, =_i, =_p\}$ |
| $\{0_{ip}, 4_{ip}\}$ | $\{3_{ip}, 4_{ip}, 03_{ii}, =_i, =_p, <\}$ |
| $\{0_{ip}, 24_{ii}\}$ | $\{14_{ii}, 24_{ii}, 03_{ii}, 34_{ii}, 04_{ii}, 44_{ii}, =_i, =_p, <\}$ |
| $\{0_{ip}, 03_{ii}\}$ | |
| $\{0_{ip}, 34_{ii}\}$ | |
| $\{0_{ip}, 04_{ii}\}$ | |
| $\{0_{ip}, 44_{ii}\}$ | |
| $\{1_{ip}, 2_{ip}\}$ | |
| $\{1_{ip}, 3_{ip}\}$ | |
| $\{1_{ip}, 4_{ip}\}$ | |
| $\{1_{ip}, 24_{ii}\}$ | |
| $\{1_{ip}, 03_{ii}\}$ | |
| $\{1_{ip}, 34_{ii}\}$ | |
| $\{1_{ip}, 04_{ii}\}$ | |
| $\{1_{ip}, 44_{ii}\}$ | |
| $\{2_{ip}, 3_{ip}\}$ | |
| $\{2_{ip}, 4_{ip}\}$ | |
| $\{2_{ip}, <\}$ | |
| $\{3_{ip}, 14_{ii}\}$ | |
| $\{3_{ip}, 24_{ii}\}$ | |
| $\{3_{ip}, 34_{ii}\}$ | |
| $\{3_{ip}, 04_{ii}\}$ | |
| $\{3_{ip}, 44_{ii}\}$ | |
| $\{4_{ip}, 14_{ii}\}$ | |
| $\{4_{ip}, 24_{ii}\}$ | |
| $\{4_{ip}, 34_{ii}\}$ | |
| $\{4_{ip}, 04_{ii}\}$ | |
| $\{4_{ip}, 44_{ii}\}$ | |

Table 17: Minimally \mathfrak{R}^+ -complete and maximally \mathfrak{R}^+ -incomplete sets. - Class: Den.

subsets of relations, and, with minimal effort, also all expressively different subsets of relations for the interesting sub-languages of interval relations only or mixed relations only. Two out of four interesting classes of linearly ordered sets are treated in Part I of this paper, while the remaining two are dealt with in the present one (Part II). There are several aspects of temporal reasoning in computer science to which this extensive study can be related:

- first-order logic over linear orders extended with temporal relations between points, intervals and mixed, is the very foundation of modal logics for temporal reasoning, and it is necessary to have a complete understanding of the former in order to deal with the latter. Indeed these first-order languages and the second-order languages based on them are the *correspondence languages* (over models and frames, respectively), for interval-based temporal logics. These logics have as yet a very imperfectly understood correspondence theory (see e.g.[VB01] and [CGP14]), and a proper understanding of the families of correspondence languages and their relative expressivity as is developed in this paper, can be seen as an important step

| \mathfrak{R}^+ | |
|--------------------------------|---|
| mcs | MIS |
| $\{0_{ip}, 2_{ip}\}$ | $\{0_{ip}, 1_{ip}, 14_{ii}, =_i, =_p, <\}$ |
| $\{0_{ip}, 3_{ip}\}$ | $\{2_{ip}, 24_{ii}, =_i, =_p, <\}$ |
| $\{0_{ip}, 4_{ip}\}$ | $\{2_{ip}, 04_{ii}, =_i, =_p\}$ |
| $\{0_{ip}, 24_{ii}\}$ | $\{3_{ip}, 4_{ip}, 03_{ii}, =_i, =_p, <\}$ |
| $\{0_{ip}, 03_{ii}\}$ | $\{14_{ii}, 24_{ii}, 03_{ii}, 34_{ii}, 04_{ii}, 44_{ii}, =_i, =_p, <\}$ |
| $\{0_{ip}, 34_{ii}\}$ | |
| $\{0_{ip}, 04_{ii}\}$ | |
| $\{0_{ip}, 44_{ii}\}$ | |
| $\{1_{ip}, 2_{ip}\}$ | |
| $\{1_{ip}, 3_{ip}\}$ | |
| $\{1_{ip}, 4_{ip}\}$ | |
| $\{1_{ip}, 24_{ii}\}$ | |
| $\{1_{ip}, 03_{ii}\}$ | |
| $\{1_{ip}, 34_{ii}\}$ | |
| $\{1_{ip}, 04_{ii}\}$ | |
| $\{1_{ip}, 44_{ii}\}$ | |
| $\{2_{ip}, 3_{ip}\}$ | |
| $\{2_{ip}, 4_{ip}\}$ | |
| $\{2_{ip}, 14_{ii}\}$ | |
| $\{2_{ip}, 24_{ii}, 04_{ii}\}$ | |
| $\{2_{ip}, 03_{ii}\}$ | |
| $\{2_{ip}, 34_{ii}\}$ | |
| $\{2_{ip}, 04_{ii}, <\}$ | |
| $\{2_{ip}, 44_{ii}\}$ | |
| $\{3_{ip}, 14_{ii}\}$ | |
| $\{3_{ip}, 24_{ii}\}$ | |
| $\{3_{ip}, 34_{ii}\}$ | |
| $\{3_{ip}, 04_{ii}\}$ | |
| $\{3_{ip}, 44_{ii}\}$ | |
| $\{4_{ip}, 14_{ii}\}$ | |
| $\{4_{ip}, 24_{ii}\}$ | |
| $\{4_{ip}, 34_{ii}\}$ | |
| $\{4_{ip}, 04_{ii}\}$ | |
| $\{4_{ip}, 44_{ii}\}$ | |

Table 18: Minimally \mathfrak{R}^+ -complete and maximally \mathfrak{R}^+ -incomplete sets. - Class: Unb.

towards the development of such a theory. As a simple example, consider the following: because of the greater expressivity of first-order languages, there is more inter-definability between sets of temporal relations in first-order languages than between the corresponding sets of modalities in propositional temporal languages. This allows one to establish connections between propositional temporal languages which would not have been directly available via the study of the languages themselves;

- automated reasoning techniques for interval-based modal logics are at their first stages; an uncommon, but promising approach is to treat them as pure modal logics

- over particular Kripke-frames, whose first-order properties are, in fact, representation theorems such as those (indirectly) treated in this paper. As a future work, we also plan to systematically study the area of representation theorems;
- the decidability of pure first-order theories extended with interval relations is well-known [Lad]; nevertheless, these results hinge on the decidability of $\text{MFO}[<]$, while we believe that they could be refined both algorithmically and computationally;
 - the study of other related languages, important in artificial intelligence, can benefit from our results, such as first-order and modal logics for spatial reasoning where basic objects are, for example, rectangles.

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